

Implementation of a Fault-Tolerant Real-time Event Channel

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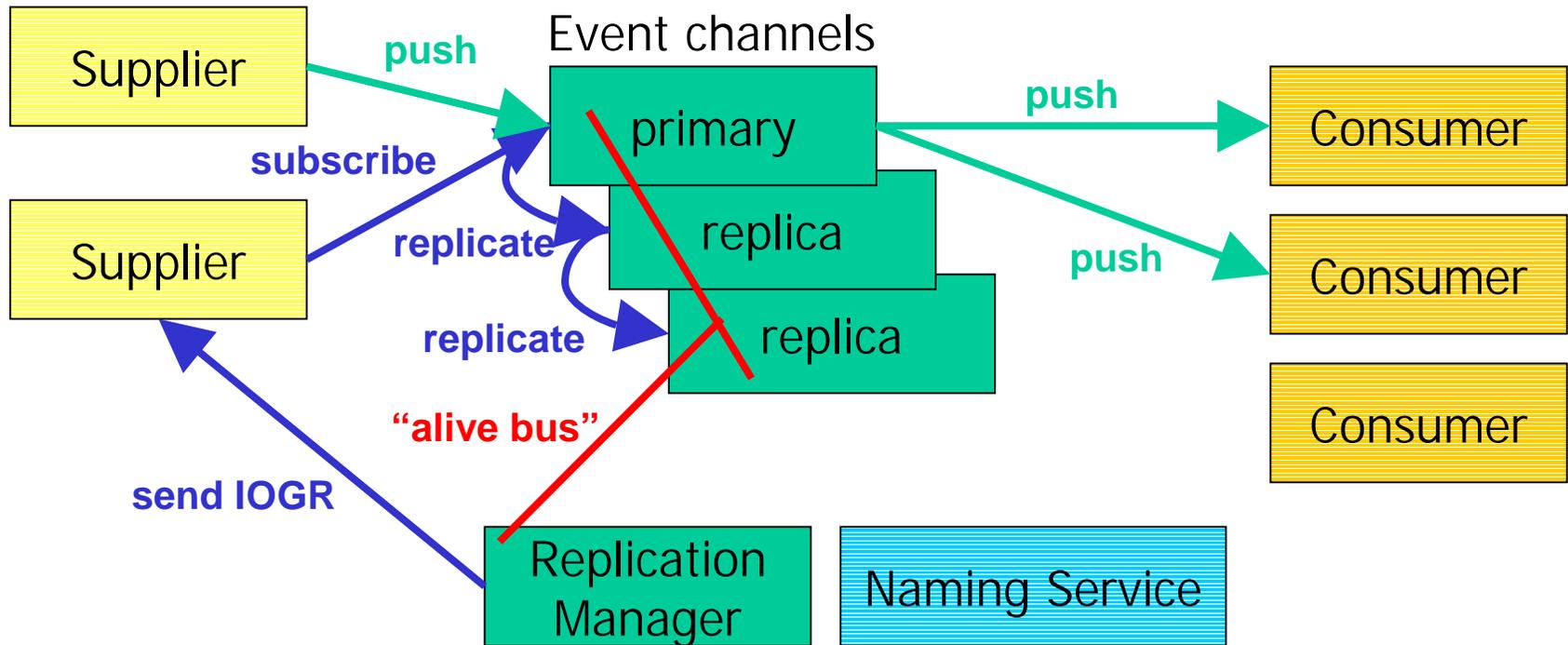
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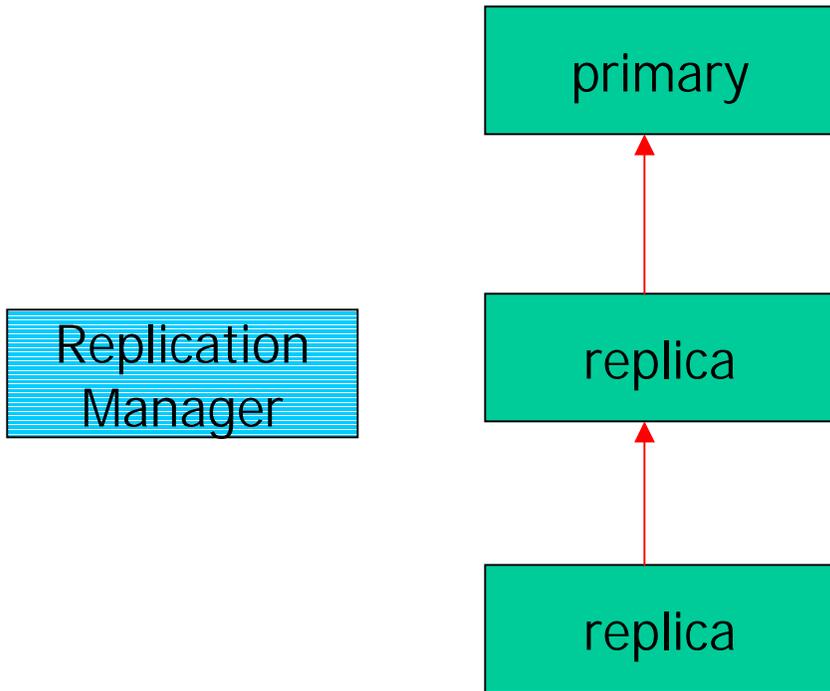
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FT/RT Event Channel (FTRTEC)



- Provide fault-tolerance (fail-stop) within real-time constraints
- Offer useful configuration knobs, *e.g.*, to Quality Connectors
 - Replicas: where and how many, transactional replication depths
- Service-level implementation
 - Less dependent on ORB-level FT features (mainly need IOGR abstraction)

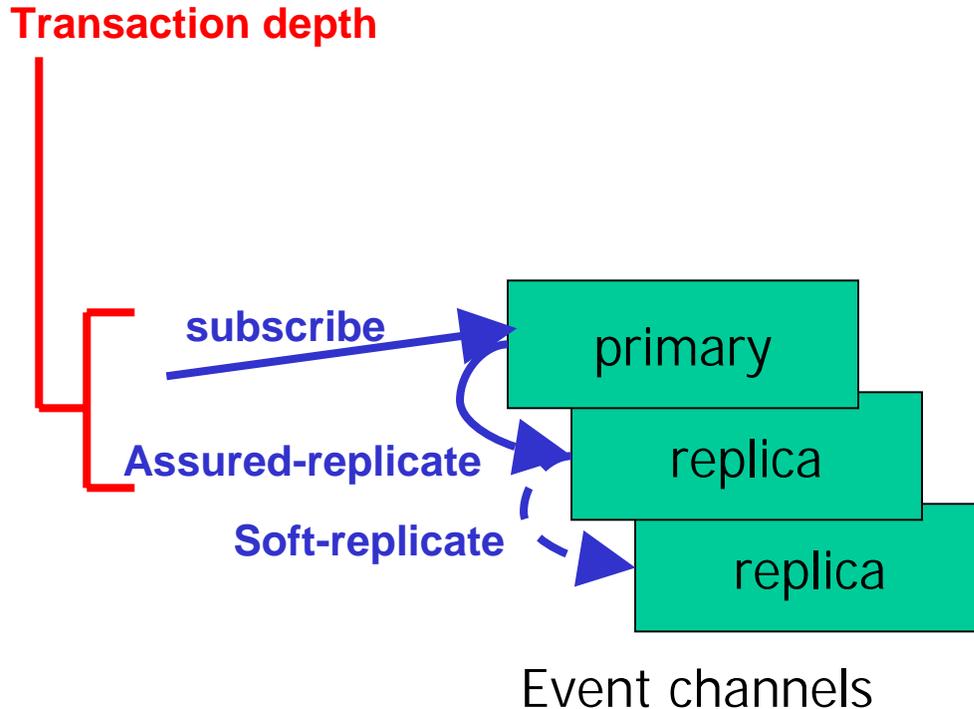
FTRTEC Fault-Detection and Fail-Over



- Maintain connections
 - Connected == alive
 - Among primary/replicas
 - To replication manager service
 - Current implementation uses TCP
- Communication of replica updates
 - Must be transactional
 - To given depth of replication
 - But need not be synchronous
 - CORBA AMI or two-way calls ok
 - But not (unreliable) one-ways

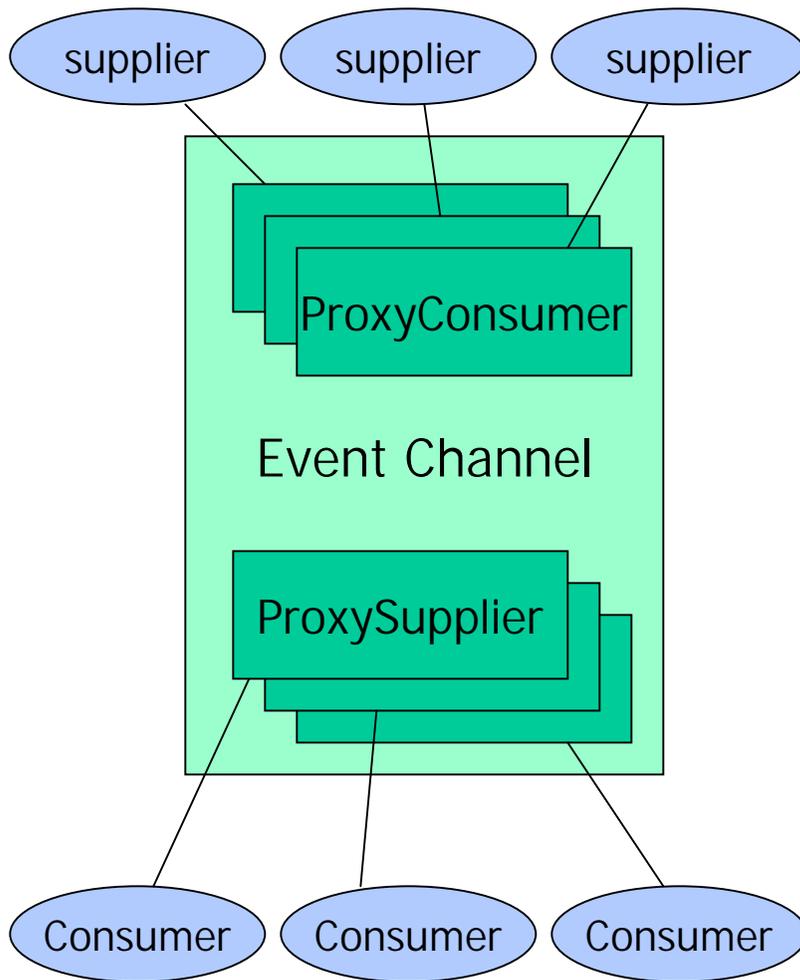
Replicated Event Channels

Subscription Replication Trade-Offs



- Risk vs. Blocking times
 - Transaction depth to tradeoff reliability and responsiveness
 - Requires two phase protocol for all replicated objects
 - Use two-way or AMI for assured-replication
 - To specified depth
 - Use oneway operations for soft-replication
 - Beyond specified depth

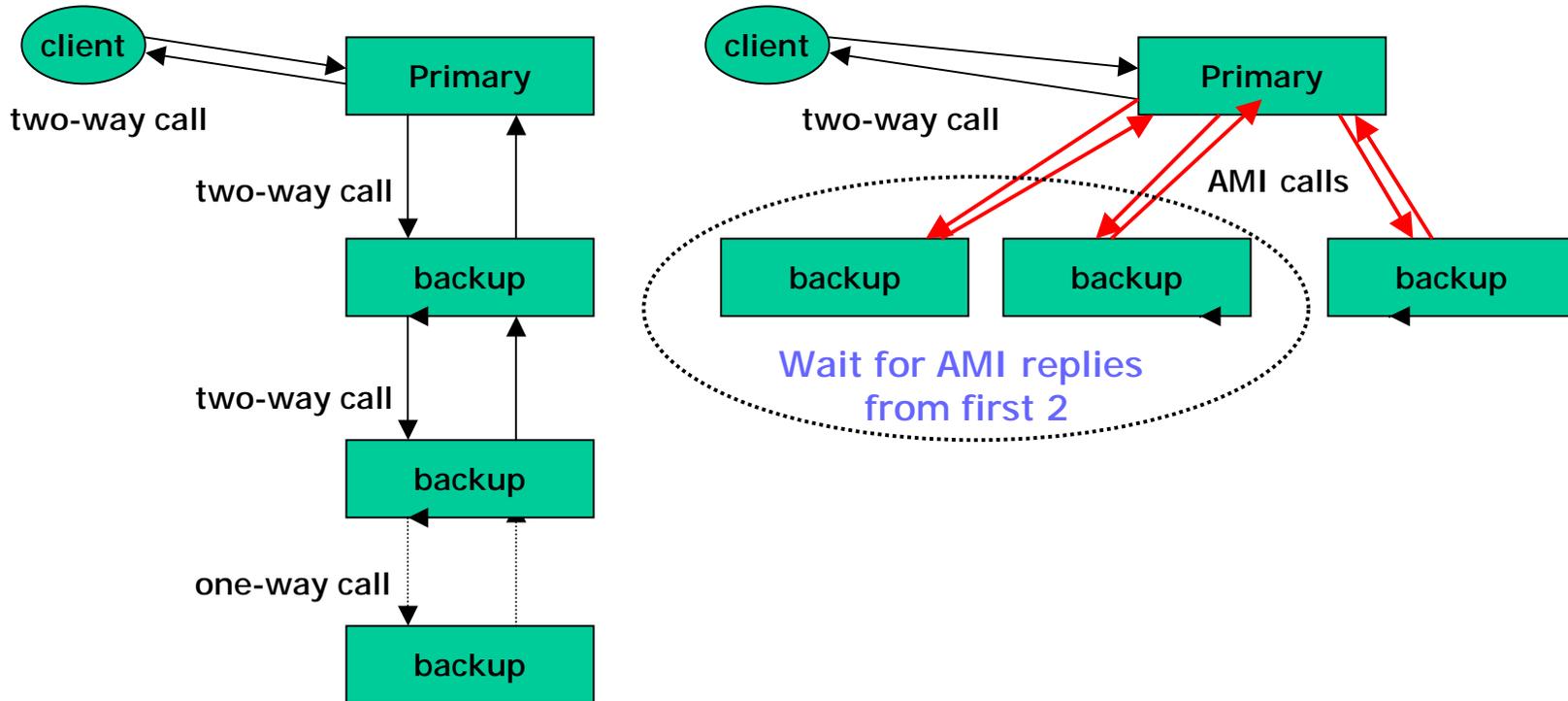
Message vs. State Replication



- Multi-faceted Objects:
 - Facade to flatten representation
 - Message-based replication
 - At object, not ORB level
- Transient vs. persistent state
 - Only replicate subscriptions currently
 - Use transactions for assurance
 - Protects the event *stream*
 - *Even during subscription*
 - Protecting events is plausible
 - Using redundant paths
 - But would require split/join semantics
 - Potential area of future work

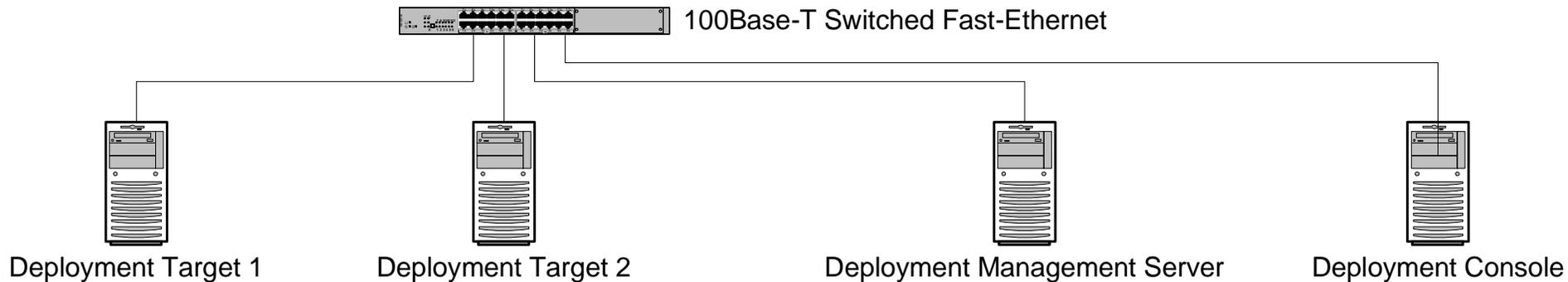
Two-Way vs. AMI Calls for Replication

Transaction depth = 3



- Initial two-way call from client to primary
- Two-way calls to the next replica result in sequential processing of updates
- AMI calls to replicas from primary instead allow overlapped processing of updates
- Offers replication speed-up for transaction depth > 2

Experimental Testbed

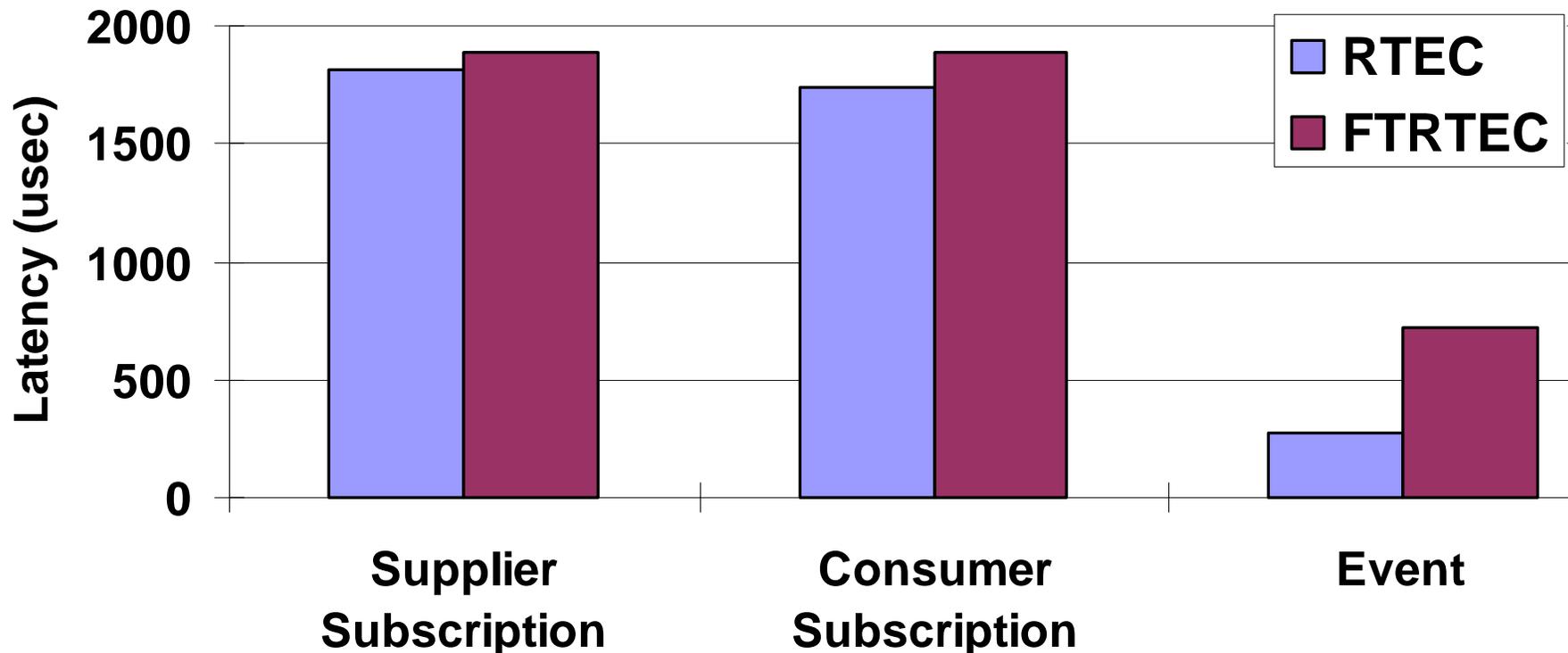


- 2 Pentium-IV 2.5 GHz machines with 500MB RAM, 512KB cache
- 2 Pentium-IV 2.8 GHz machines with 500MB RAM, 512KB cache
- KURT-Linux 2.4.18
- 100 Base-T Ethernet, isolated network for experiment runs
- ACE version 5.3.5 / TAO version 1.3.5 (pre-release version)
- Experiments were run as root in real-time scheduling class

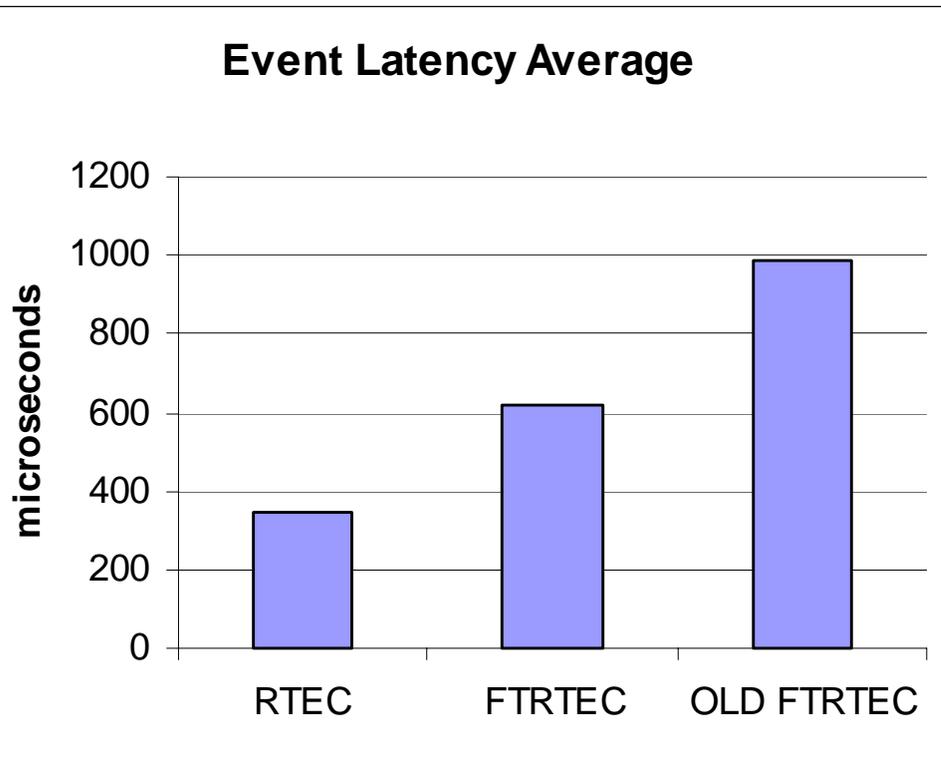
Original FTRTEC vs. RTEC Benchmarks

- Metrics: event throughput, subscription latency (TAO RTEC is baseline)
- TAO FT/RT EC vs. TAO RTEC, no replicas for FTRTEC
- 2 2.8 GHz Pentium boxes, KURT-Linux, Ethernet
- Suppliers/consumers on same machine, EC on a separate machine
- FTRTEC two-way event overhead/fail-over trade-off (AMI optimization WIP)

Baseline EC vs. FTRTEC Latency Comparison



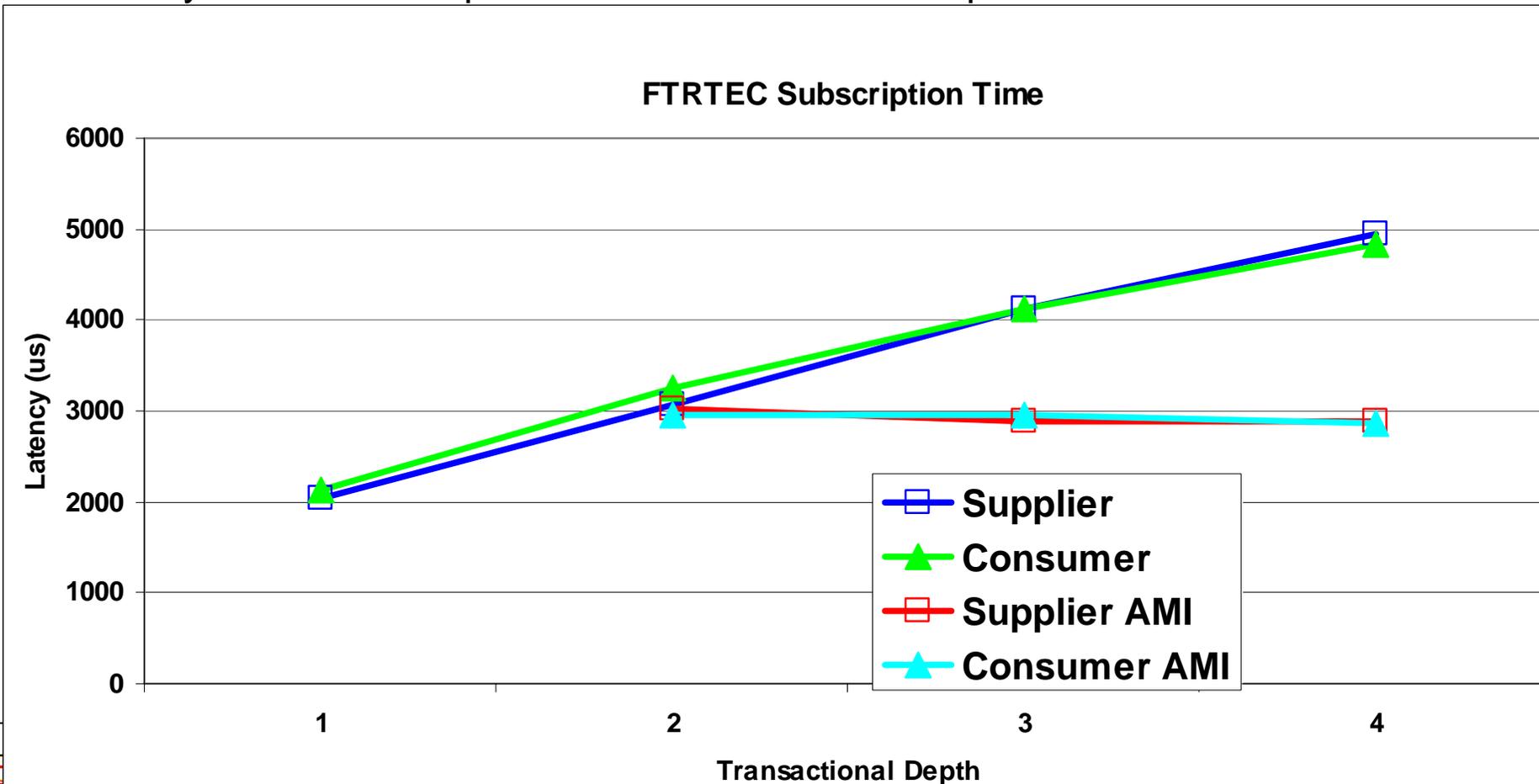
Current FTRTEC vs. RTEC Latency Benchmark



- We made several key FTRTEC improvements
 - Client EC gateway uses a separate ORB (event pushes bypass IOGR interceptors)
 - Moved IOGR service context processing to not impact events
 - Removed a redundant service context id check from FTCORBA core
- And re-ran experiments with pre-release TAO 1.4

Subscription Time Scalability

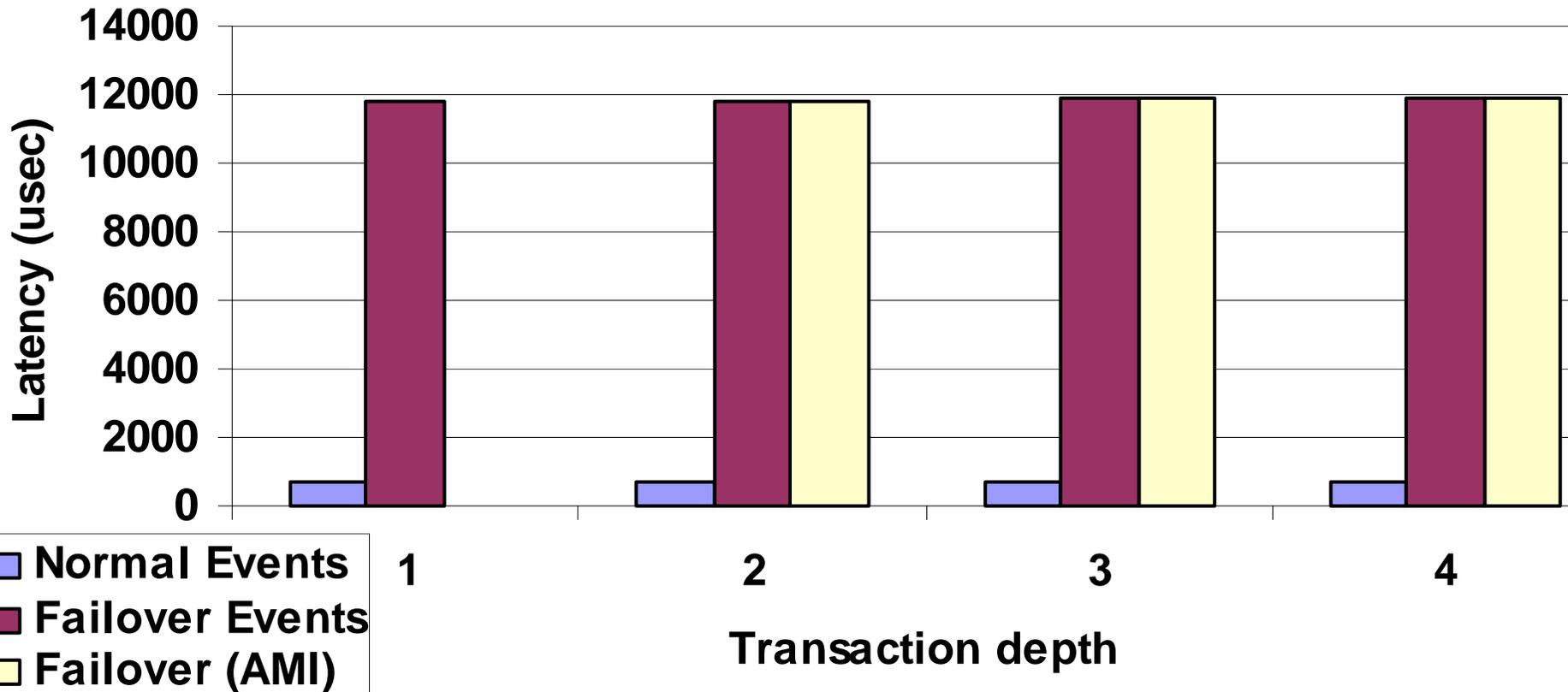
- Metrics: subscription latency/depth with two-way vs. AMI replication
- 2 2.8 GHz + 2 2.6 GHz Pentium boxes, KURT-Linux, Ethernet
- AMI/two-way active replication from primary to replicas up to transaction depth
- One-way semi-active replication after transaction depth



FTRTEC Normal, Fail-Over Event Latency

- Metrics: fail-over latency versus normal event latency
- 2 2.8 GHz & 2 2.6 GHz Pentium boxes, KURT-Linux, Ethernet
- Approximately a factor of 15 greater latency for fail-over (still, > 80 Hz)

FTRTEC Normal and Failover Event Latency



Concluding Remarks

- Using CORBA AMI optimizes primary to replica updates
 - Exploits overlapped processing of updates
 - While still allowing transactional semantics
- Able to trade-off real-time performance with fault-tolerance
- Removing replication mechanisms from event paths is key
- FT/RT Event Channel is available open-source
 - ACE version 5.3.5 / TAO version 1.3.5 and later
- FT/RT EC integrated into LMCO MINERS demo in December
- Thanks to Joe Cross (DARPA), and Sylvester Fernandez and David Mattioli (LMCO)
- Thanks also to Andy Gokhale, Bala Natarajan, and Doug Schmidt at Vanderbilt University