



**Component-based approach
for real-time and embedded
systems**

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<http://www.ist-compare.org>

COMPonent Approach for Real-time and Embedded

- ❑ Collaborative european IST project (running until end 2006)
- ❑ Focus : *component-based approach for real-time and embedded systems*
- ❑ Project consortium:



THALES

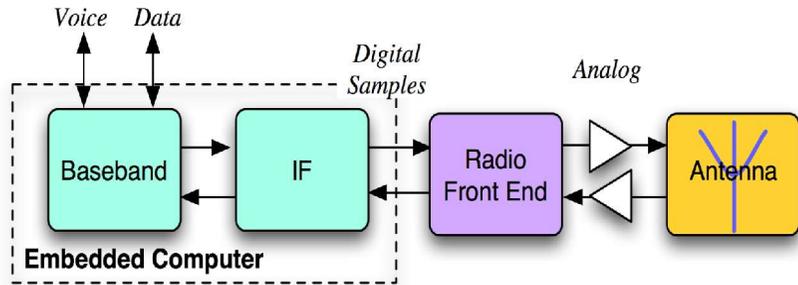


 PRISMTECH
The Integration Server Company

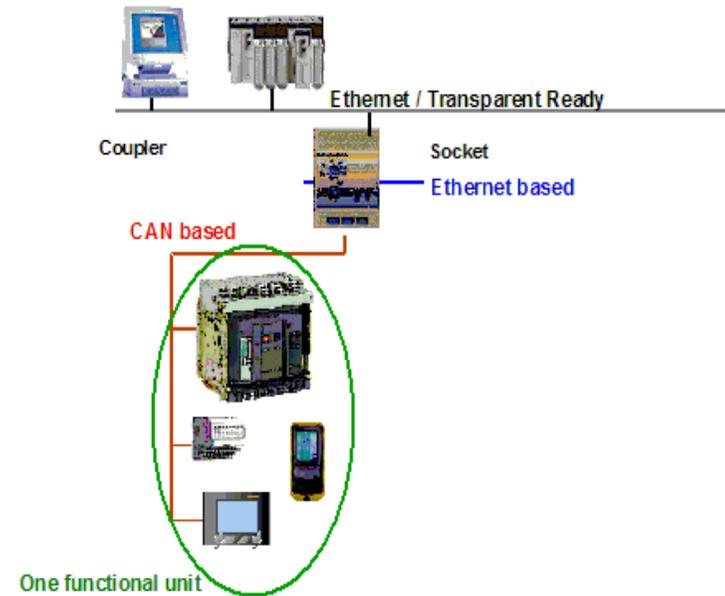
TRIALOG

Schneider
 Electric

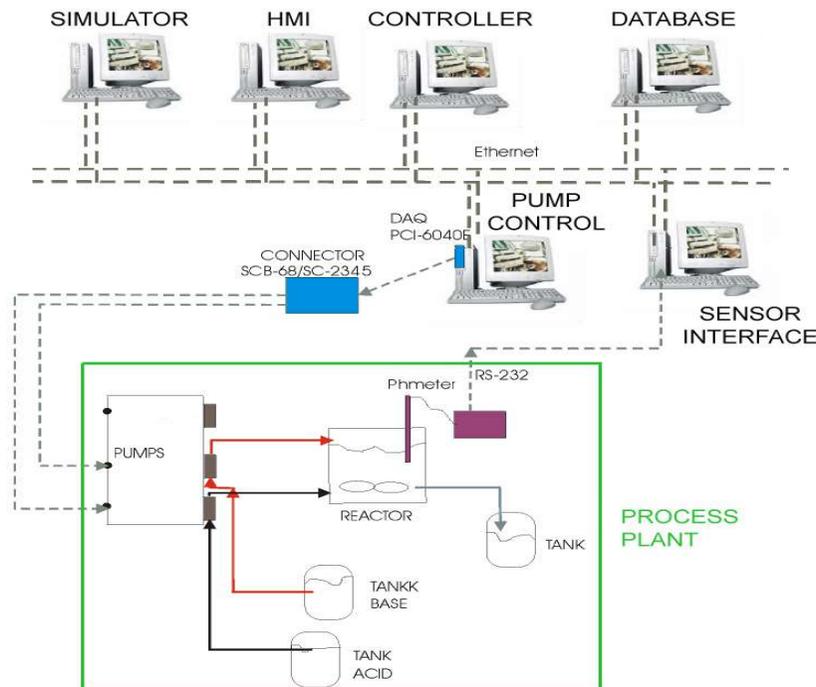
- Software defined radio



- Electrical distribution



- distributed control systems



+ automotive systems

□ RT/E systems sw-engineering via component approach

⇒ Motivations

- Designing real-time systems like communication ones is hard
 - **real-time** constraints (schedule, synchronization), **embedded** (code adapted to hardware target **platform** ?)
 - hard **integration** times: increased complexity, more functionalities, more variability
 - **openness** constraints (e.g Software defined radio)
- Reusing previous developments in slightly different scopes can still be difficult
 - example : portability on different RTOSes
- Reuse (when made possible) granularity is traditionally too coarse

⇒ Relevance of Component based approach:

- fosters **reusability** of system functions in the scope of product lines and beyond
- Resulting system sound architecture, eases analysis of system
- **portability**: allocation of components to different kind of resources without effort
- Offers a pattern for a clean **separation of concerns**

⇒ **But** today, no broadly applicable component approach for real-time and embedded systems is available

What we propose

⇒ New approach for sw architecture of real-time embedded systems

□ Definition of a component model and accompanying toolset

→ following a smartly layered approach (it's a framework)

→ enforcing separation of concerns

★ e.g real-time scheduling / algorithmic code

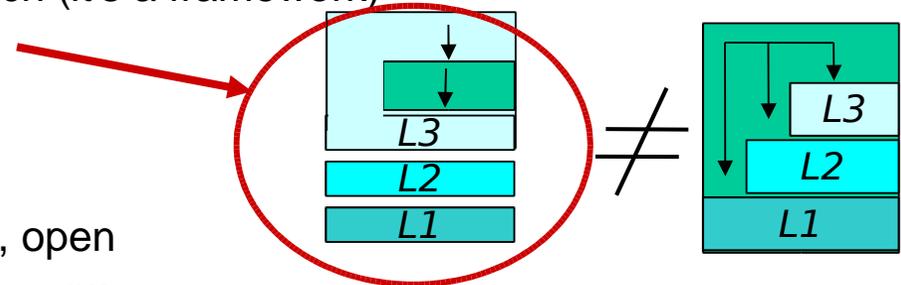
→ Lightweight, versatile, value-added, open

★ reuse component across (truly) different execution environments

★ ability to embed nothing more than what needed in the infrastructure

★ coming with a collection of pre-existing services and architectural patterns

⇒ Builds on / adapts / profiles existing standards (OMG Lw-CCM)



□ This talk concentrates on:

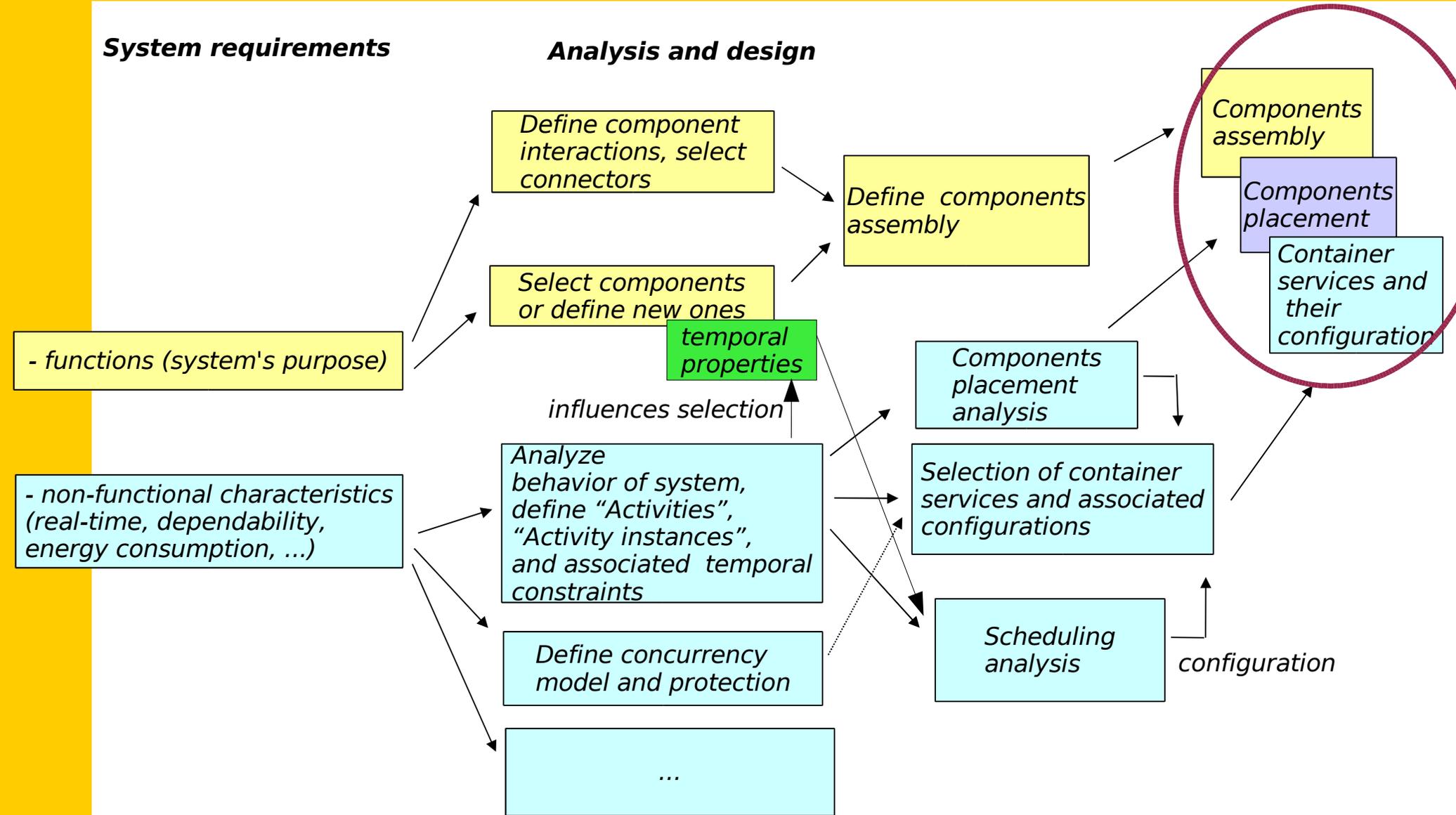
⇒ open container supporting real-time features

⇒ reusable interactions support

- ⇒ Follow and extend **standard** approaches
 - OMG lightweight CCM, OMG D&C specification, QoS for CCM
 - push the obtained extensions for standardization
- ⇒ **Separation of concerns**
 - ultimately reusable software versus embedding it into a variety of non-functional environments (real-time platforms, etc ...)
- ⇒ Deal with **real-time**
 - integrate basic real-time ingredients, scheduling strategies choices, concurrency management and realization transparently to components code
- ⇒ Allow smooth integration of / gatewaying to **legacy** code / systems
 - crucial for acceptance in an industrial context
- ⇒ Examine carefully **tradeoffs** in terms of **tailoring** / profiling standards as well as implementation techniques to get optimal performance and footprint
- ⇒ Realize **implementations** of the fwk with following underlying RT platforms:
 - real-time CORBA (e*ORB) on VxWorks / PPC and OSE ck / TI C55 DSP (SDR)
 - on OSEK (automotive domain static RTOS) + OSEK Com (Electrical breaker)

System requirements

Analysis and design

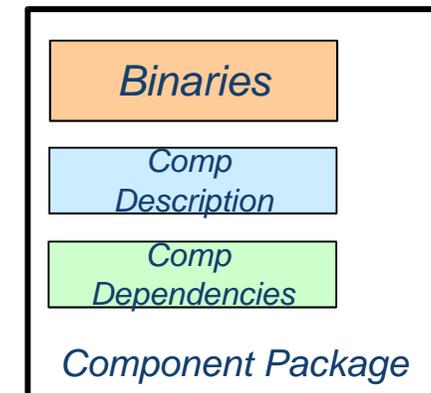
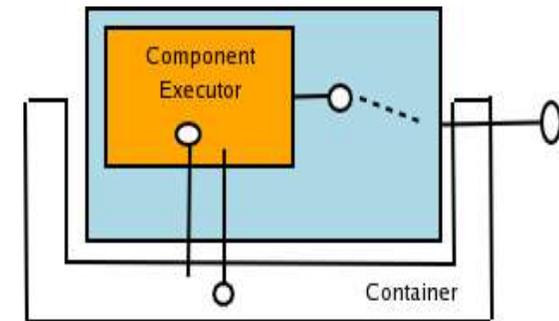
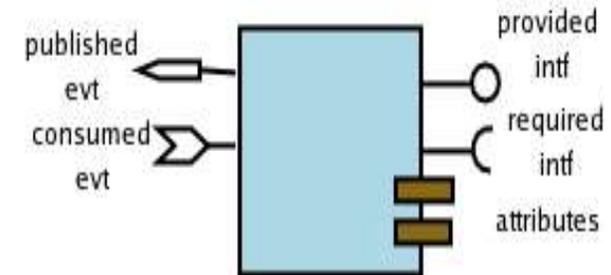


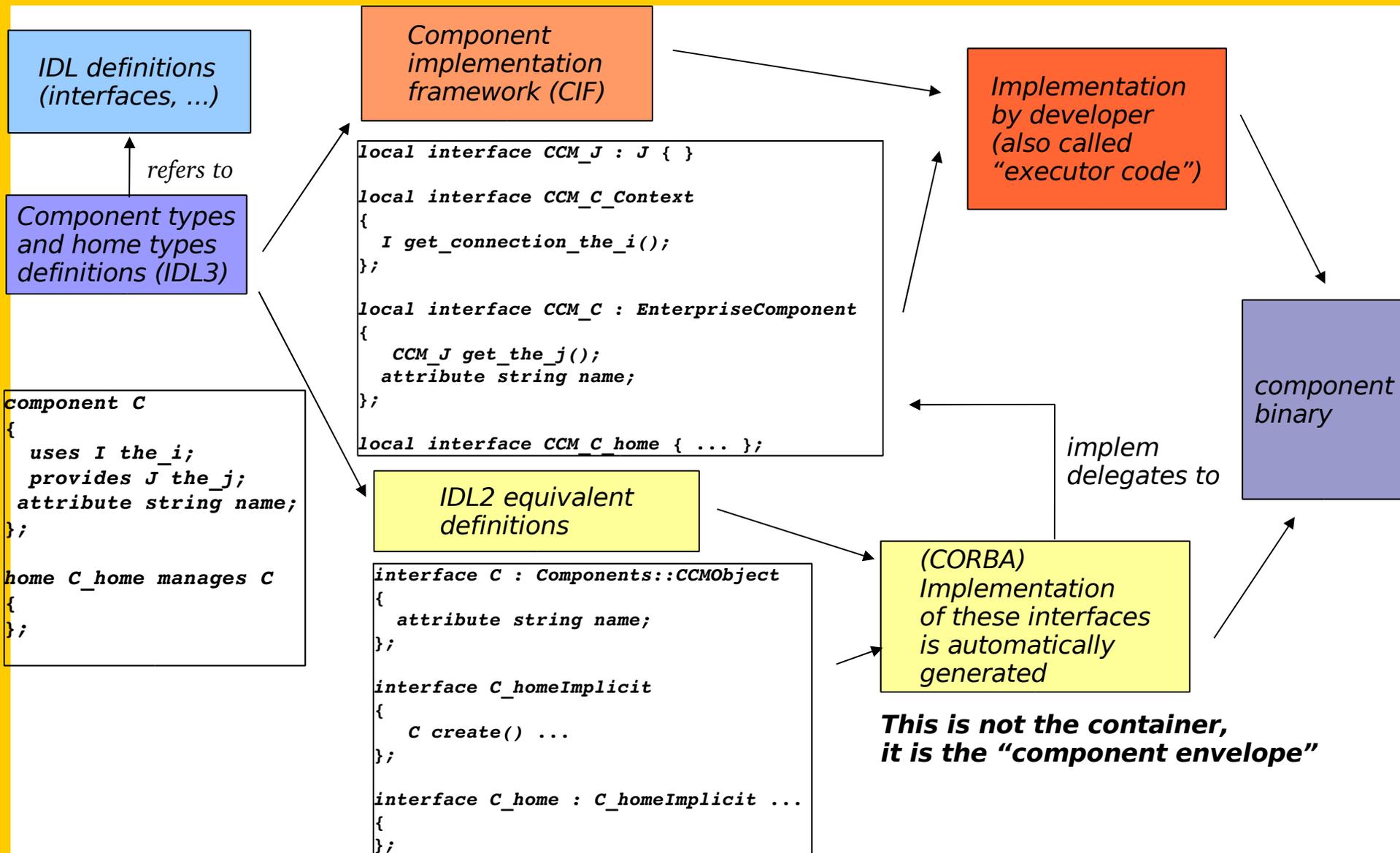
- A component model
 - ⇒ description of component types
 - provided / required services
 - events published / consumed

- An execution environment
 - ⇒ containers
 - shield components developers from middleware technical concerns
 - Underlying CORBA middleware support
 - but not only (in our extensions)

- A packaging model
 - ⇒ software-package format
 - ⇒ self-descriptive : contents, dependencies

- A deployment infrastructure
 - ⇒ assembly descriptors + runtime support





Extension of CCM capabilities in terms of interactions

⇒ Currently :

- Synchronous method invocation
- (1,n) push-push typed events propagation

Without any QoS

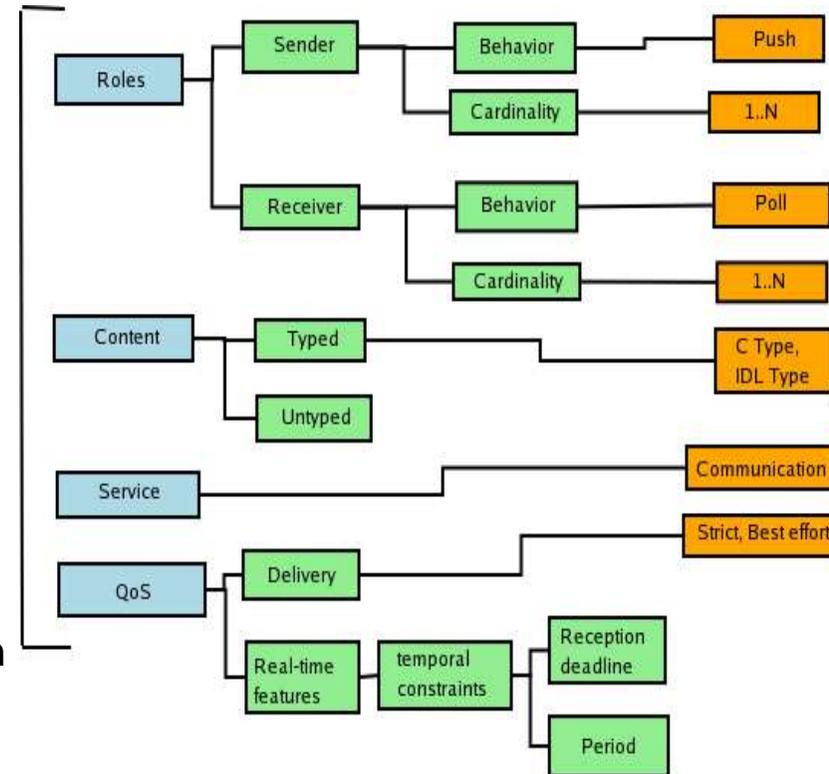
⇒ Other architectural styles of interest in real-time systems

- Pipes and filters
 - classical in communication protocols domain
- Variants of publish-subscribe
 - push/pull, typed / untyped events
- Blackboard pattern

⇒ Characterization of an interaction via

- Participants to the interaction
 - ★ roles, and cardinalities
- Type of service
 - ★ communication or/and synchronization
- QoS features
 - ★ synchronicity, dynamic behaviour

example



- Introduced
 - ⇒ the Connector concept as in the ADLs to abstract interactions
 - ⇒ Mechanisms to provide Connector implementations
 - ability to define new port types (in addition to : provides, uses, consumes, ...)
 - ability to define new connector types
 - rules for connectors implementation and packaging
 - ⇒ Extensions of application descriptions
 - ⇒ Extensions to infrastructure to deal with connectors
- We have a set of classical connectors available:
 - Deferred synchronous calls (callback based, poll based)
 - Event bus mechanisms
 - Optimized direct method calls
- We can provide easily(w/o modifying the obtained comp. model):
 - blackboard connector
 - connectors for ad-hoc messaging scheme / ad-hoc transport (i.e non CORBA based)
 - streaming connector

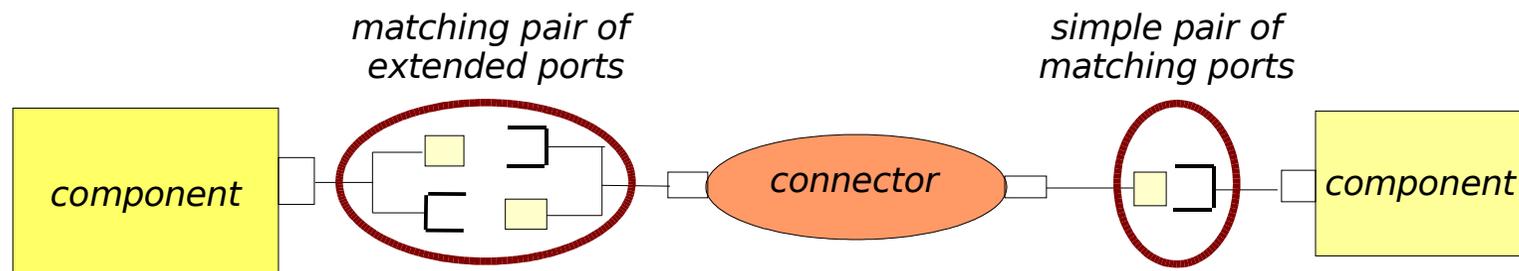
Component / Connector model :

⇒ “Extended port concept”:

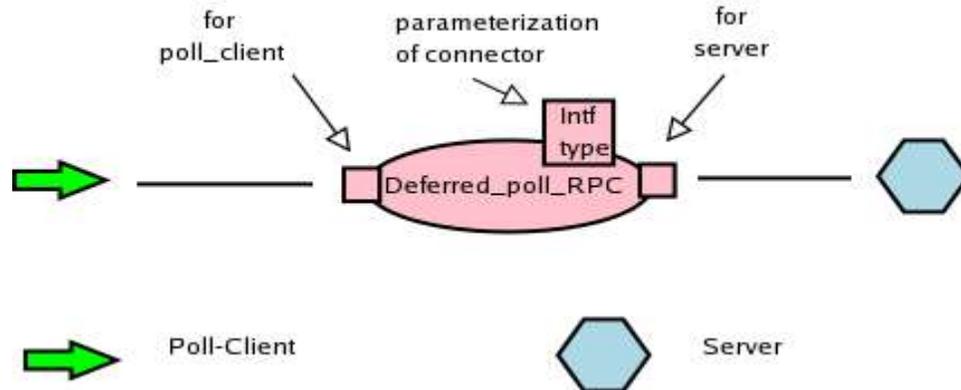
- can define the port type needed in case of specific interaction model.
- can be parameterized by an interface type or event type (like C++ templates).
- extended port as collection of needed / used interfaces + associated semantics
 - Component model becomes a little bit closer to UML2 component model

⇒ Connector concept :

- stands for interaction entity having some “extended ports” as well as attributes
 - looks like components but see some slides later for differences



⇒ Deferred synchronous call with polling



```

Port_type RPC_Poll_Client<I>
{
  provides trsf(I) client_port;
};

connector RPC_Poll<I> {
  provides_port RPC_Poll_Client<I> fpc;
  uses I for_server;
};

```

```

component Client_component {
  uses_port RPC_Poll_Client<I> my_port;
};

```

```

component Server_component {
  provides I a_port;
};

```

```

interface I {
  long foo( in string s );
};

local interface long_Poller {
  bool is_available();
  void get_result( out long result );
  void check_for_exception( ... );
};

// interface effectively used by caller
// (refers to trsf(I) above)

```

```

local interface RPC_Poll_I {
  long_Poller foo( in string s );
};

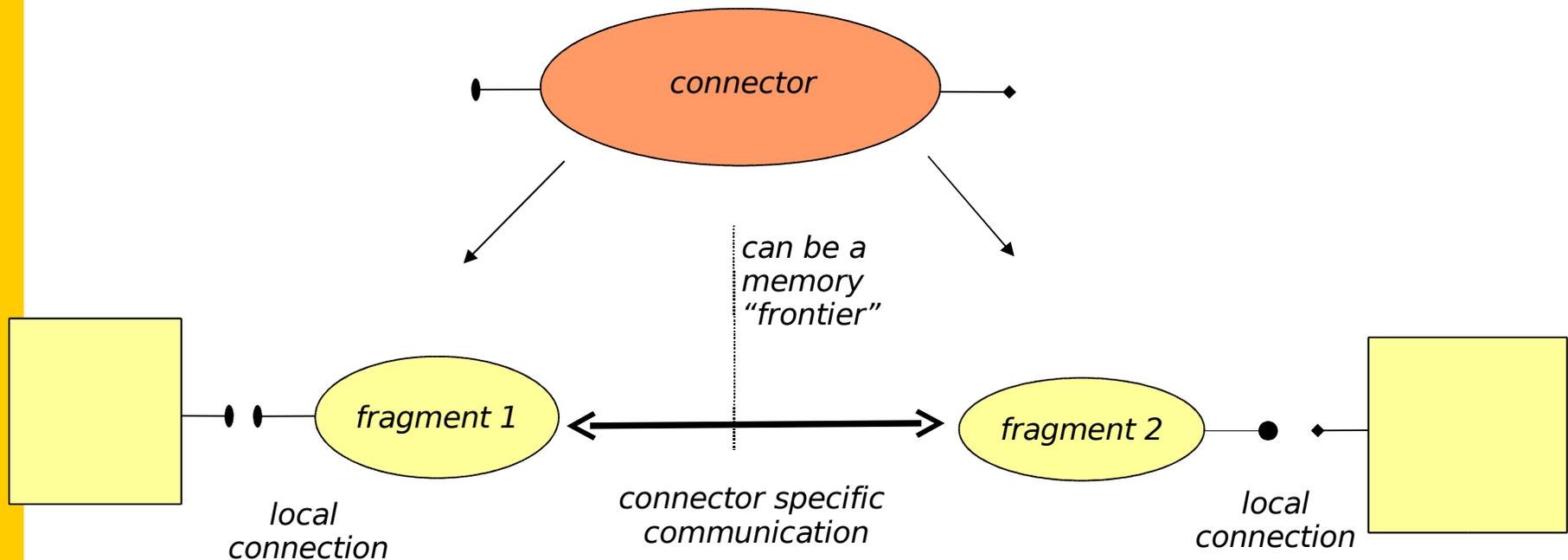
```

⇒ Client uses an interface derived from I by transformation:

- namely **RPC_Poll_I**
 - call `foo()` passing arg,
 - does not block
 - obtains “long_Poller” to check for result later

Connectors implementation :

- ⇒ is naturally “fragmented” (connector parts co-located with components involved in interaction)
- ⇒ connector fragments implementations follow the component approach
 - they implement the “IDL3 extended ports” concept



Definition of an open (extensible) container to support various combinations of **plug-ins** for **non-functional** aspects handling.

⇒ Motivations to provide a container:

- Provide true **separation of concerns**

- Means in practice: code no more polluted with tasks creations, synchronization, ...

- ★ i.e. these aspects are to be integrated transparently to component code and described (not hard-coded).

- This is what Containers (in component technologies) are meant for!

- Allow creation of an **open** and **customizable** execution environment via containers

- To fit exactly application needs (in terms of threading, communication, etc.)

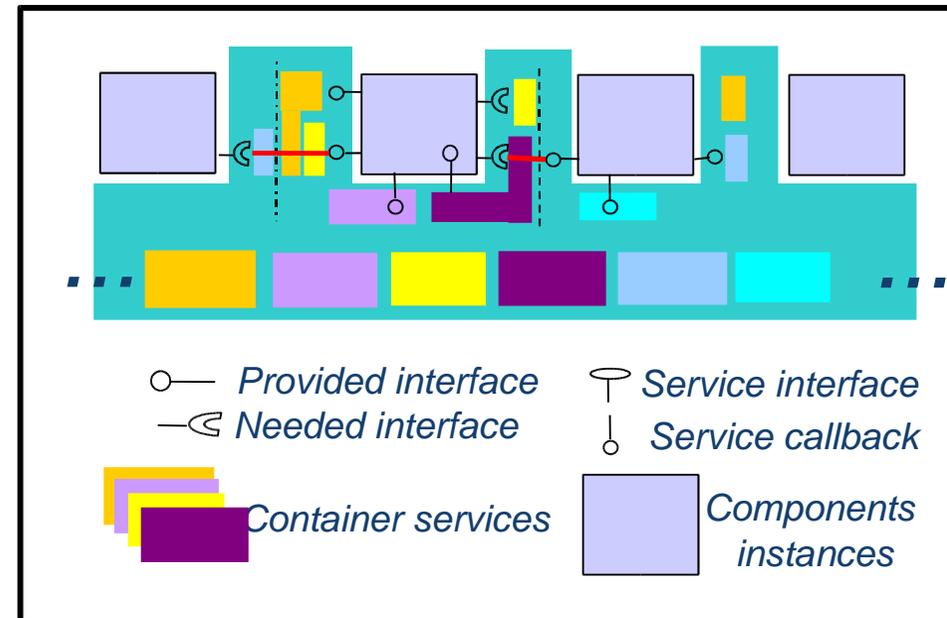
- Coming with a set of pre-existing services (e.g scheduling, thread-pools, memory pools, ...)

⇒ Approach:

- Build custom containers starting from modules: called **container services**

- Define generic interfaces to implement container-integrable modules

- Container as host for:
 - homes, components
 - service plug-ins: are aggregation of objects involved at well-specified occasions (corresponding to “integration points” below) in the container
- provides what we call “integration points” for services:
 - requests *interception*
 - control at *components creation time*
 - control at *components connection time*
 - integration of *connector fragments*
- APIs (1) are defined for insertion of service at “integration points”
 - see next slides for details
- API defined for Service plug-ins
 - Service plug-ins install objects implementing interfaces from (1) at “service integration points”



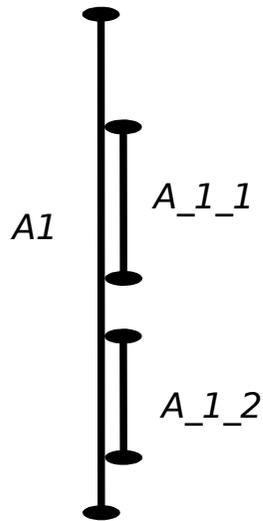
Integration points:

<i>Pre and Post - request interception</i>	<p>interceptor :</p> <p>bound to comp instance + method (smart proxies techniques)</p> <p>access to request (end-to-end) context</p> <p>other granularities possible</p>
<i>Request interception – In place of</i>	Corresponds to integration of Connector fragments
<i>components creation time and servants activation / deactivation time</i>	activating the component servants on a POA (having specific policies) and in a specific manner
<i>Container internal interfaces (component context)</i>	for container services interfaces needed to be presented to component
<i>Components Callbacks</i>	interfaces implemented by component that the container uses

Notion of “Activity” as a design concept:

- ⇒ builds on the notion of “distributable” thread of RT-CORBA 2
- ⇒ Activity : logical entity corresponding to a behavior of the system
 - can be described as a sequence of actions whose starting point can be identified, as well as stop point (which can be implicit)
 - associated to a component instance, a facet, and a method
- ⇒ Activity instance :
 - realization of an Activity at particular (absolute) time, under a particular context.
- ⇒ Real-time formal vocabulary
 - Classical concepts:
 - release time, relative deadline, jitter, importance
 - as annotation of activity instances

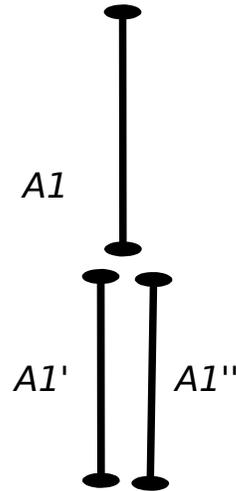
Activity compositions patterns (can be combined)



nesting of Activities

A1, A_1_1, and A_1_2 can be annotated with RT-constraints

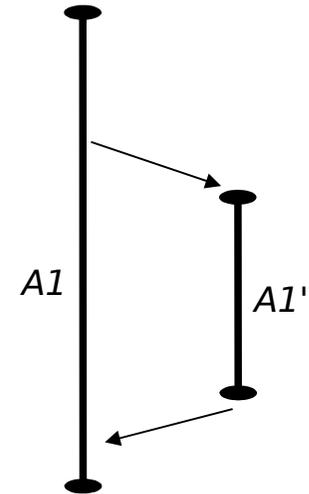
Typically : situation arises in the case of nested component calls



chained transfer of activity

*A1' and A1'' start points is A1 end point.
Release time of A1' and A1''
= end time of A1*

*Typically : situation arises as natural precedence constraints,
A1 not triggering directly A1' and A1''*



**fork and join
or
fork and forget**

A1' starts at specific component interaction

- ❑ Container services identified and considered:
 - ⇒ realize the activity patterns from previous slide
 - ❑ activity creation,
 - timers, interaction with component wrapping hardware
 - ❑ nested activity creation
 - specific components interaction (as part as “specific” activities)
 - ❑ new (forked) activity
 - same
 - ❑ transfer of activity

- ❑ Ingredients / container mechanisms to deal with “Activities”
 - ❑ end-to-end request context (carrying “Activity context” tailored for app)
 - ❑ Pre and post interception
 - ❑ combined with synchronous, deferred synchronous, async connectors

- Real-time CORBA 1.0 :
 - ⇒ is a priority based real-time middleware
- Schedulability analysis outputs:
 - ⇒ set of priorities for
 - activities, nested activities, forked activities, and transferred activities
- Container services identified and involved:
 - ⇒ RT-CORBA Client propagated policy application
 - usage of an “On activation” integration point
 - ⇒ Application of the set of priorities:
 - pre or post – interceptors (coming with Activity mgt plugins from previous slide):
 - configured with mapping : activity context to priority
 - using the RT-CORBA API to set priorities (i.e RTCORBA::Current)

Mapping can also be done onto CORBA dynamic scheduling and deadline schedulers.

Can also be done onto other mechanisms like typically direct POSIX primitives

Mechanisms to automate management of concurrency:

- ⇒ Intent: manage as far as possible concurrent accesses protection without presenting Lock APIs to the component developer.
- ⇒ Some ideas and considered techniques
 - Lock-based technique :
 - lock protections at component instance level, facet level, ...
 - done by container plug-in : lock / unlock with pre / post interception
 - placement of locks is not easy, too coarse granularity
 - ★ but description of app activities can help to make the analysis.
 - Alternatives (under investigation):
 - implement “Transactional software memory” patterns
 - ★ can be automated to some extent
 - directly use “lock-free” data structures when possible
 - ★ responsibility put on component developer.

- ❑ Assessment of the container and real-time techniques considered
- ❑ Implementation of framework:
 - ⇒ onto e*ORB SDR C++ onto VxWorks / PowerPC
 - ⇒ onto e*ORB SDR C onto OSE compact kernel on TI C55x DSP
 - ⇒ minimalistic implementation onto OSEK / automotive in EC++
 - ❑ need to optimize very much and allow a fully static approach
- ❑ Implementation of test beds to validate the whole approach:
 - ⇒ Software defined radio waveform (THALES)
 - ⇒ Electrical breaker (Schneider-Electric)

- ❑ More information and details at:
⇒ <http://www.ist-compare.org>

- ❑ Thank you for your attention