Implementation of a
Fault-Tolerant Real-time Event Channel

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FT/RT Event Channel (FTRTEC)

- Provide fault-tolerance (fail-stop) within real-time constraints
- Offer useful configuration knobs, e.g., to Quality Connectors
  - Replicas: where and how many, transactional replication depths
- Service-level implementation
  - Less dependent on ORB-level FT features (mainly need IOGR abstraction)
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FTRTEC Fault-Detection and Fail-Over

- Maintain connections
  - Connected == alive
  - Among primary/replicas
  - To replication manager service
  - Current implementation uses TCP
- Communication of replica updates
  - Must be transactional
    - To given depth of replication
  - But need not be synchronous
    - CORBA AMI or two-way calls ok
    - But not (unreliable) one-ways

Replicated Event Channels
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Subscription Replication Trade-Offs

- Risk vs. Blocking times
  - Transaction depth to tradeoff reliability and responsiveness
  - Requires two phase protocol for all replicated objects
  - Use two-way or AMI for assured-replication
    - To specified depth
  - Use oneway operations for soft-replication
    - Beyond specified depth
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Message vs. State Replication

- Multi-faceted Objects:
  - Facade to flatten representation
  - Message-based replication
    - At object, not ORB level
- Transient vs. persistent state
  - Only replicate subscriptions currently
    - Use transactions for assurance
    - Protects the event stream
    - Even during subscription
  - Protecting events is plausible
    - Using redundant paths
    - But would require split/join semantics
    - Potential area of future work
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Two-Way vs. AMI Calls for Replication

Transaction depth = 3

- Initial two-way call from client to primary
- Two-way calls to the next replica result in sequential processing of updates
- AMI calls to replicas from primary instead allow overlapped processing of updates
- Offers replication speed-up for transaction depth > 2
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Experimental Testbed

- 2 Pentium-IV 2.5 GHz machines with 500MB RAM, 512KB cache
- 2 Pentium-IV 2.8 GHz machines with 500MB RAM, 512KB cache
- KURT-Linux 2.4.18
- 100 Base-T Ethernet, isolated network for experiment runs
- ACE version 5.3.5 / TAO version 1.3.5 (pre-release version)
- Experiments were run as root in real-time scheduling class
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Original FTRTEC vs. RTEC Benchmarks

- Metrics: event throughput, subscription latency (TAO RTEC is baseline)
- TAO FT/RT EC vs. TAO RTEC, no replicas for FTRTEC
- 2 2.8 GHz Pentium boxes, KURT-Linux, Ethernet
- Suppliers/consumers on same machine, EC on a separate machine
- FTRTEC two-way event overhead/fail-over trade-off (AMI optimization WIP)

Baseline EC vs. FTRTEC Latency Comparison
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Current FTRTEC vs. RTEC Latency Benchmark

- We made several key FTRTEC improvements
  - Client EC gateway uses a separate ORB (event pushes bypass IOGR interceptors)
  - Moved IOGR service context processing to not impact events
  - Removed a redundant service context id check from FTCORBA core
- And re-ran experiments with pre-release TAO 1.4
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Subscription Time Scalability

- Metrics: subscription latency/depth with two-way vs. AMI replication
- 2 2.8 GHz + 2 2.6 GHz Pentium boxes, KURT-Linux, Ethernet
- AMI/two-way active replication from primary to replicas up to transaction depth
- One-way semi-active replication after transaction depth

![Graph showing FTRTEC Subscription Time](image-url)
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**FTRTEC Normal, Fail-Over Event Latency**

- Metrics: fail-over latency versus normal event latency
- 2 2.8 GHz & 2 2.6 GHz Pentium boxes, KURT-Linux, Ethernet
- Approximately a factor of 15 greater latency for fail-over (still, > 80 Hz)

### FTRTEC Normal and Failover Event Latency

![Graph showing comparison of normal and failover event latency with transaction depth](image-url)
Concluding Remarks

• Using CORBA AMI optimizes primary to replica updates
  – Exploits overlapped processing of updates
  – While still allowing transactional semantics
• Able to trade-off real-time performance with fault-tolerance
• Removing replication mechanisms from event paths is key
• FT/RT Event Channel is available open-source
  – ACE version 5.3.5 / TAO version 1.3.5 and later
• FT/RT EC integrated into LMCO MINERS demo in December
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