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# Robotic Technology Component Specification

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## **Preface**

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Times/Times New Roman - 10 pt.: Standard body text

Helvetica/Arial - 10 pt. Bold: OMG Interface Definition Language (OMG IDL) and syntax elements.

Courier - 10 pt. Bold: Programming language elements.

Helvetica/Arial - 10 pt: Exceptions

**Note** – Terms that appear in *italics* are defined in the glossary. Italic text also represents the name of a document, specification, or other publication.

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## 1 Scope

## 1.1 Overview

This document defines a component model and certain important infrastructure services applicable to the domain of robotics software development. It includes a Platform-Independent Model (PIM) expressed in UML and Platform-Specific Models (PSMs) expressed in OMG IDL.

The sophistication of robotic and other intelligent systems has increased in recent years, allowing such systems to be applied to an ever-increasing range of problems. Such systems are no longer constrained to the factory floor, but are now used in a variety of commercial and non-commercial applications. They have even begun to enter our homes. The technologies used in these applications, which can be applied not only to standalone robots but also to ubiquitous computing and other more intelligent electrical devices, we call "Robotic Technology" (RT).

As the level of structural and behavioral complexity in RT systems has increased, so too has the need for technologies supporting the integration of those systems. A systematic methodology and infrastructure for the local and distributed composition of modular functional components is needed.

This document specifies a component model that meets these requirements of RT systems. We refer to a component that supports the integration of RT systems as a Robotic Technology Component (RTC) [RTC RFP]. An RTC is a logical representation of a hardware and/or software entity that provides well-known functionality and services. By extending the general-purpose component functionality of UML with direct support for domain-specific structural and behavioral design patterns, RTCs can serve as powerful building blocks in an RT system. Developers can combine RTCs from multiple vendors into a single application, allowing them to create more flexible designs more quickly than before.

This specification does not seek to replace existing component models (e.g., UML components, Lightweight CCM, Software Radio components). Instead, it focuses on those structural and behavioral features required by RT applications that are not addressed by pre- existing models. It is anticipated that some implementers of this specification will choose to implement it in terms of another, more general, component model.

## 1.2 Platform-Independent Model

The PIM described herein is applicable to a very wide range of robotics and controls applications, from fine-grained local components to coarser-grained components communicating over a network. Its focus is at a somewhat higher level of abstraction and domain-specificity than that of more general purpose models (e.g., [CCM]) and it includes direct support for several important design patterns that are central to the architecture of many robotics applications.

The PIM consists of three parts:

- Lightweight RTC. A simple model containing definitions of concepts such as component, port, and the like. This section provides the foundation on which the subsequent sections build. See Section 5.2.
- Execution semantics. Extensions to Lightweight RTC to directly support critical design patterns used in robotics applications such as periodic sampled data processing, discrete event/stimulus response processing, modes of operation, etc. See Section 5.3.
- Introspection. An API allowing for the examination of components, ports, connections, etc. at runtime. See Section 5.4.

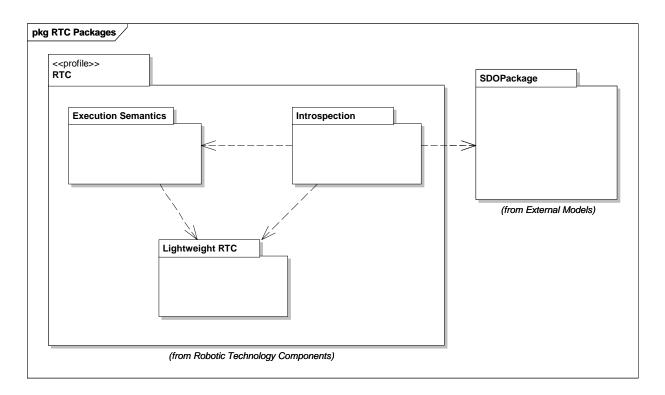


Figure 1.1 - RTC packages

Figure 1.1 shows the relationships among the sections described above.

## 1.3 Platform-Specific Models

Three PSMs are described:

- 1. Local. Components reside on the same network node and communicate over direct object references without the mediation of a network or network-centric middleware such as CORBA.
- 2. Lightweight CCM [CCM]. Most components are assumed to be distributed relative to one another; they communicate using a CCM-based middleware.
- 3. CORBA [CORBA]. Most components are assumed to be distributed relative to one another; they communicate using a CORBA-based middleware.

The Local PSM is primarily applicable to relatively fine-grained components executing within a single application. It offers the potential for very low latency, high determinism communications (on the order of a small number of function calls) and is thus appropriate for components or applications with hard real-time requirements.

The [CORBA]- and [CCM]-based PSMs are primarily applicable when most components are distributed over a network. These components are likely to be coarser grained, more autonomous, and have higher tolerances for latency, jitter, and communications failure than components that communicate locally. Example applications include groups of service robots communicating within a home and mobile phones exchanging voice calls or other data within a cellular network.

## 2 Conformance and Compliance

Support for Lightweight RTC (see Section 5.2) is fundamental and obligatory for all implementations.

In addition, implementations are strongly urged to support at least one of the following optional conformance points, as they provide greater domain-specific value to robotics applications than Lightweight RTC alone.

- Periodic Sampled Data Processing (see Section 5.3.1)
- Stimulus Response Processing (see Section 5.3.2)
- Modes (see Section 5.3.3)
- Introspection (see Section 5.4)

At least one of the following PSMs must be implemented for each of the conformance points in the list above to which conformance is claimed.

- Local communications (see Section 6.3)
- [CCM]-based communications (see Section 6.4)
- [CORBA]-based communications (see Section 6.5)

Interoperability across programming languages is beyond the scope of the Local PSM. Conformance is therefore relative to a particular IDL-to-programming language mapping. Providers must state the language mapping(s) for which their implementation is conformant. The following is an example statement of conformance:

*Middleware X* is conformant to the Lightweight RTC, Periodic Sampled Data Processing, and Introspection conformance points of the Robotic Technology Component (RTC) Specification on the Local platform in the C++ language, as defined by the OMG-standard IDL-to-C++ mapping, version 1.1 (formal/03-06-03).

## 3 References

## 3.1 Normative References

[CORBA] Common Object Request Broker Architecture (CORBA) Core Specification, version 3.0.3

Specification: formal/2004-03-12

http://www.omg.org/technology/documents/formal/corba\_iiop.htm

#### [CCM] CORBA Component Model, version 4.0

http://www.omg.org/technology/documents/formal/components.htm

#### [RTC RFP] RFP for Robot Technology Components

http://www.omg.org/cgi-bin/doc?ptc/2005-09-01

### [SDO] Super Distributed Objects

http://www.omg.org/cgi-bin/doc?formal/2004-11-01

#### [UML] Unified Modeling Language, Superstructure Specification version 2.1.1

http://www.omg.org/technology/documents/formal/uml.htm

## 3.2 Non-normative References

#### [CCM STREAM] Streams for CCM Specification

This specification is still in finalization as of this writing. The document number is ptc/05-07-01. http://www.omg.org/cgi-bin/doc?ptc/2005-07-01

## [CCM UML] UML Profile for CORBA Components, version 1.0

Specification: formal/05-07-06 Supporting files: formal/05-08-03

http://www.omg.org/technology/documents/formal/profile\_ccm.htm

## [CORBA UML] UML Profile for CORBA Specification, version 1.0

Specification: formal/2002-04-01

http://www.omg.org/technology/documents/formal/profile\_corba.htm

## [DC] Deployment and Configuration of Component-based Distributed Applications, version 1.0

http://www.omg.org/cgi-bin/doc?ptc/2004-08-02

## [DDS] Data Distribution Service, version 1.1

http://www.omg.org/cgi-bin/doc?formal/05-12-04

#### [RT CORBA] RealTime-CORBA Specification, version 2.0

Specification: formal/2003-11-01

http://www.omg.org/technology/documents/formal/RT\_dynamic.htm

#### [SWRADIO] PIM and PSM for Software Radio Components

http://www.omg.org/cgi-bin/doc?dtc/2005-09-05

## 4 Additional Information

## 4.1 Requirements

RT applications typically require the following features:

• Components of varying granularity: Software must be hierarchically decomposable to an arbitrary depth. A component model must therefore support components of varying granularity. A simple algorithm performing some mathematical or control function may be encapsulated as a fine-grained component. On the other extreme, an entire humanoid robot may be represented as a single autonomous, coarse-grained component. A component developer must be free to choose the granularity appropriate for a given component and to compose components of differing granularity within a single application.

- Autonomous and passive components: Some components must operate autonomously and in parallel with one another. Such components must own an independent task of execution in which to carry out their work (e.g., real-time feedback control). In other cases, however, separate threads for each component is an unnecessary expense (e.g., returning a value from a method invocation). The mapping of components to threads of execution must therefore be flexible.
- Real-time capabilities: Some tasks within an RT application (such as servo control) require hard real-time capabilities, including very low latencies and a high degree of determinism.
- Systematic execution: The execution ordering of modules within an application frequently follows one of a few idiomatic design patterns, such as sampled execution or finite state machines. The component model should support such design patterns.
- Software reuse: Components must be strongly encapsulated for maximum reusability.
- Network and communication independence: Different applications have different requirements with respect to intercomponent communication. A component model must therefore be independent of any particular communications medium or protocol.

## 4.2 Relationships to Existing OMG Specifications

## 4.2.1 Platform-Independent Model

The PIM is defined in UML 2.1.1 [UML].

The full RTC specification extends Lightweight RTC (described in this document) and Super Distributed Objects (described in [SDO]).

## 4.2.2 Platform-Specific Models

The PSMs are defined in part using OMG IDL. The Lightweight CCM PSM depends on that specification [CCM].

Note that the transformation of the PIM into CCM does not rely on the UML Profile for CORBA [CORBA UML] or the UML Profile for CORBA Components [CCM UML].

- The purpose of the CCM PSM in this specification is to describe models based on UML components using CCM.
- The purpose of [CORBA UML] and [CCM UML] is the opposite: to describe CORBA- and CCM-based applications
  using UML. They do not use UML components or interfaces (IDL components and interfaces are represented as UML
  classes).

Therefore, this specification maps the relevant concepts to CCM directly. More information may be found in Chapter 6.

## 4.3 Acknowledgements

The following companies submitted and/or supported parts of this specification.

- Japan Robot Association (JARA)
- National Institute of Advanced Industrial Science and Technology (AIST)
- Real-Time Innovations (RTI)

- Seoul National University (SNU)
- Technologic Arts Incorporated

## 5 Platform Independent Model

## 5.1 Format and Conventions

Classes are described in this PIM using tables of the following format. A more detailed description of each member follows the table.

<class name=""></class>					
attributes	attributes				
<attribute name=""> <attribute type=""></attribute></attribute>					
operations					
<pre><operation name=""></operation></pre>		<return type=""></return>			
	<pre><parameter name=""></parameter></pre>	<pre><parameter type=""></parameter></pre>			

Enumerated constants are modeled as attributes of the enumeration type.

Operation parameters can contain a modifier "in," "out," or "inout" ahead of the parameter name. If the modifier is omitted, it is implied that the parameter is an "in" parameter.

In some cases the operation parameters and/or return value(s) are a collection with elements of a given <type>. This scenario is indicated with the notation "<type> []." This notation does not imply that it will be implemented as an array. The actual implementation is defined by the PSM; it may end up being mapped to a sequence, an array, or some other kind of collection.

For example, the class named MyClass below has a single attribute named my\_attribute of type long and a single operation my\_operation that returns a long. The operation takes four parameters. The first, param1, is an output parameter of type long; the second, param2, an input-output parameter of type long; the third, param3, is an input parameter (the "in" modifier is implied by omission) of type long; and the fourth, param4, is also an input parameter of type collection of longs.

MyClass		
attributes		
my_attribute long		
operations		
my_operation		long
	out: param1	long

inout: param2	long
param3	long
in: param4	long[]

## 5.2 Lightweight RTC

The lightweight RT component model defines interfaces and stereotype to adapt UML components to the RT domain. The UML 2.0 Components and Composite Structures packages—and those packages on which they depend, directly or indirectly—are considered normative. A non-normative simplified version of that model is shown in Figure 5.1.

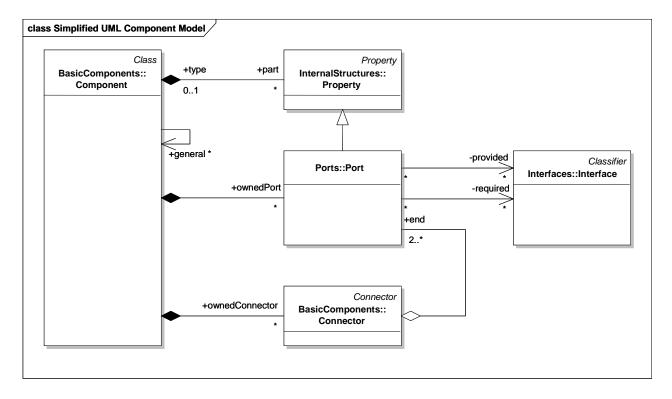


Figure 5.1 - Simplified depiction of components and their constituents from [UML]

## Components

From [UML]:

A component is a self contained unit that encapsulates the state and behavior of a number of classifiers. A component specifies a formal contract of the services that it provides to its clients and those that it requires from other components or services in the system in terms of its provided and required interfaces.

A component is a substitutable unit that can be replaced at design time or run-time by a component that offers equivalent functionality based on compatibility of its interfaces. As long as the environment obeys the constraints expressed by the provided and required interfaces of a component, it will be able to interact with this environment. Similarly, a system can be extended by adding new component types that add new functionality.

#### **Ports**

#### From [UML]:

Ports represent interaction points between a classifier and its environment. The interfaces associated with a port specify the nature of the interactions that may occur over a port. The required interfaces of a port characterize the requests that may be made from the classifier to its environment through this port. The provided interfaces of a port characterize requests to the classifier that its environment may make through this port.

Note that a port may expose multiple interfaces, and those interfaces may be of mixed polarity (i.e., some may be required and others provided). This feature is important for the expression of callback semantics or any other two-way communication contract.

#### **Composite Components**

Component instances may be composed together to form new component types; these latter are called *composite component types* or sometimes *component assemblies*. The component instances that comprise the internal structure of a composite component are referred to as its *parts*. Those parts are connected to one another and to their containing composite component via *connectors* between their respective ports. Further detail is provided in [UML].

## 5.2.1 ReturnCode\_t

#### Description

A number of operations in this specification will need to report potential error conditions to their clients. This task shall be accomplished by means of operation "return codes" of type **ReturnCode** t<sup>1</sup>.

#### **Semantics**

Operations in the PIM that do not return a value of type **ReturnCode\_t** shall report errors in the following ways, depending on their return type:

- If an operation normally returns a positive numerical value (such as **get\_rate**, see Section 5.2.2.6.4), it shall indicate failure by returning a negative value.
- If an operation normally returns an object reference (such as RTObject::get\_component\_profile, see Section 5.4.2.2.1), it shall indicate failure by returning a nil reference.

#### **Attributes**

ReturnCode_t		
attributes		
OK	ReturnCode_t	
ERROR	ReturnCode_t	

<sup>1.</sup> *Implementation notes*: Depending on its target programming language and execution environment, a PSM may express ReturnCode\_t as an actual return value or as an exception. This specification does not mandate exceptions because many platforms that are important for robotics developers do not support them well. (This error reporting convention is also used by [DDS].)

BAD_PARAMETER	ReturnCode_t	
UNSUPPORTED	ReturnCode_t	
OUT_OF_RESOURCES	ReturnCode_t	
PRECONDITION_NOT_MET	ReturnCode_t	
no operations		

## 5.2.1.1 OK

## **Description**

The operation completed successfully.

## 5.2.1.2 ERROR

#### Description

The operation failed with a generic, unspecified error.

## 5.2.1.3 BAD\_PARAMETER

#### **Description**

The operation failed because an illegal argument was passed to it.

#### 5.2.1.4 UNSUPPORTED

## **Description**

The operation is unsupported by the implementation (e.g., it belongs to a compliance point that is not implemented).

## 5.2.1.5 OUT\_OF\_RESOURCES

#### **Description**

The target of the operation ran out of the resources needed to complete the operation.

## 5.2.1.6 PRECONDITION\_NOT\_MET

## **Description**

A pre-condition for the operation was not met.

## 5.2.2 Components

This section defines the stereotypes and interfaces that define conforming components. In particular, any conforming component must be extended by the stereotype lightweightRTComponent or some specialization thereof (such as rtComponent; see Section 5.4.2.1). Such a component shall be referred to as a "lightweight robotic technology component," or "lightweight RTC."

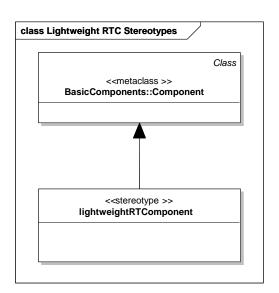


Figure 5.2 - Lightweight RTC Stereotypes

## 5.2.2.1 lightweightRTComponent

### Description

The lightweightRTComponent stereotype extends a UML component (i.e., UML::Components::BasicComponents::Component [UML]).

A component so extended conforms to the lifecycle described in Section 5.2.2.3 and allows clients to monitor and respond to changes in that lifecycle by means of the **ComponentAction** interface. The following figure is a non-normative example.

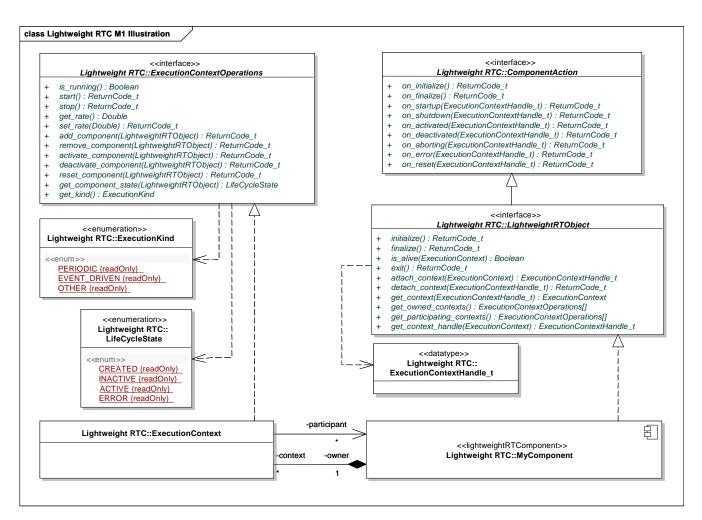


Figure 5.3 - Lightweight RTC M1 Illustration

#### **Constraints**

- A component extended by the lightweightRTComponent stereotype shall realize the LightweightRTObject interface.
- Although the UML meta-class Component allows any number of supertypes, at most one supertype may be another (lightweight) RTC.

## 5.2.2.2 LightweightRTObject

## **Description**

This interface is realized by all lightweight RTCs (as required by the **lightweightRTComponent** stereotype). It defines the states and transitions through which all RTCs will pass from the time they are created until the time they are destroyed.

#### Generalizations

• ComponentAction

#### **Semantics**

#### Initialization

An RTC begins in the Created state; at this point, it has been instantiated but not yet fully initialized. Note that this state is highly implementation-dependent. For example, it may correspond to the invocation of a constructor in languages that support that concept, but not all languages do. Furthermore, how soon this state is entered before initialize is invoked is implementation-dependent. Therefore, it should be relied on by RTC implementers only to the minimum extent possible.

An RTC that has completed its initialization and has not been finalized is said to be Alive.

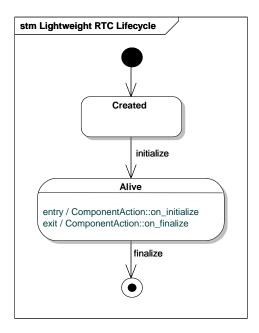


Figure 5.4 - Lightweight RTC Lifecycle

#### **Execution Context**

An RTC in the Alive state may participate in any number of execution contexts (see Section 5.2.2.5). These contexts shall be represented to an RTC as distinct instances of the **ExecutionContext** class. The **ExecutionContext** manages the behavior of each RTC that participates in it. This relationship is defined by the following state machine, which is embedded within the **ExecutionContext**'s own lifecycle (see Figure 5.5). Each participating RTC is represented as a separate parallel region.

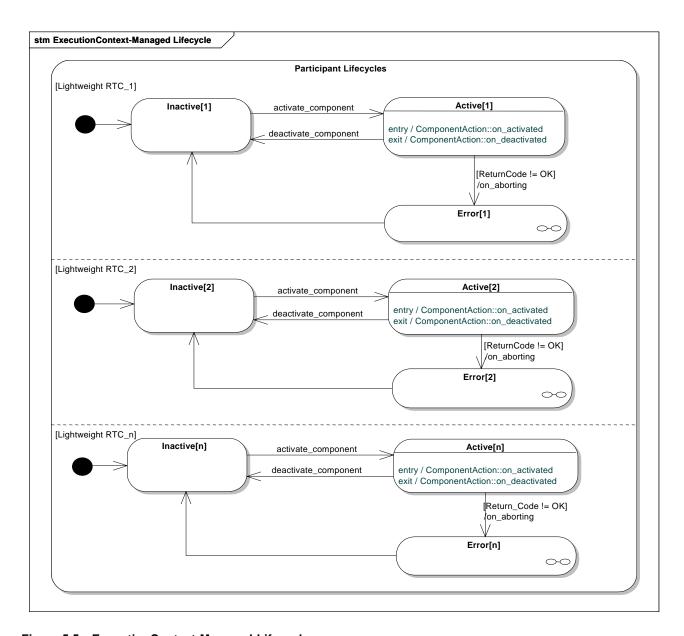


Figure 5.5 - ExecutionContext-Managed Lifecycle

Relative to a given execution context, an RTC may either be Active, Inactive, or in Error. When the RTC is Active in a Running execution context, the **ComponentAction** callbacks (see Section 5.2.2.4) shall be invoked as appropriate for the context's **ExecutionKind**. The callbacks shall not be invoked relative to that context when either the RTC is Inactive in that context or the context is Stopped. (Note that starting and stopping an execution context shall not impact whether its participating RTCs are Active or Inactive.)

It may be that a given RTC does not directly participate in any execution contexts. Such an RTC is referred to as *passive*. A passive RTC may provide services to other components upon request. At any other time, it shall not be required to perform any ongoing activity of its own; therefore, instances of such an RTC typically exist only as parts (directly or indirectly) of a containing active RTC.

## Error Handling

If an operation fails while the RTC is Active in a given execution context, the RTC will transition to the Error state corresponding to that context. While the RTC is in Error, the ComponentAction::on\_error callback will be invoked in place of those callbacks that would otherwise have been invoked according to the context's ExecutionKind. For example, if the kind is PERIODIC, on\_error shall be invoked instead of the pair of on\_execute, and on\_state\_update.

When an RTC is in Error, it may be reset. If resetting is successful, the RTC shall return to the Inactive state. If resetting is unsuccessful, it shall remain in the Error state.

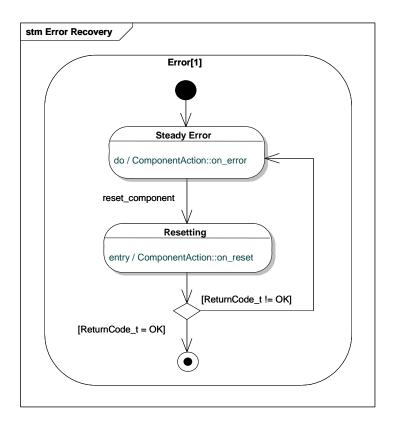


Figure 5.6 - Error Recovery

#### **Constraints**

• A class that realizes this interface must conform to the protocol state machine shown above.

## **Operations**

LightweightRTObject
no attributes
operations

initialize		ReturnCode_t
finalize		ReturnCode_t
is_alive		Boolean
	context	ExecutionContext
exit		ReturnCode_t
attach_context		ExecutionContextHandle_t
	context	ExecutionContext
detach_context		ReturnCode_t
	exec_handle	ExecutionContextHandle_t
get_context		ExecutionContext
	exec_handle	ExecutionContextHandle_t
get_owned_contexts		ExecutionContextOperations[]
get_participating_contexts		ExecutionContextOperations[]
get_context_handle		ExecutionContextHandle_t
	context	ExecutionContext

### 5.2.2.2.1 initialize

## **Description**

Initialize the RTC that realizes this interface.

## **Semantics**

The invocation of this operation shall result in the invocation of the callback ComponentAction::on initialize.

## **Constraints**

- An RTC may be initialized only while it is in the Created state. Any attempt to invoke this operation while in another state shall fail with ReturnCode t::PRECONDITION NOT MET.
- Application developers are not expected to call this operation directly; it exists for use by the RTC infrastructure.

#### 5.2.2.2.2 finalize

#### Description

Finalize the RTC that realizes this interface, preparing it for destruction.

#### **Semantics**

This invocation of this operation shall result in the invocation of the callback ComponentAction::on finalize.

#### **Constraints**

- An RTC may not be finalized while it is participating in any execution context. It must first be removed with **ExecutionContextOperations::remove\_component**. Otherwise, this operation shall fail with **ReturnCode t::PRECONDITION NOT MET**. See Figure 5.9.
- An RTC may not be finalized while it is in the Created state. Any attempt to invoke this operation while in that state shall fail with ReturnCode t::PRECONDITION NOT MET.
- Application developers are not expected to call this operation directly; it exists for use by the RTC infrastructure.

#### 5.2.2.2.3 is alive

#### Description

A component is alive or not regardless of the execution context from which it is observed. However, whether or not it is Active, Inactive, or in Error is dependent on the execution context(s) (see Figure 5.7) in which it is running. That is, it may be Active in one context but Inactive in another. Therefore, this operation shall report whether this RTC is *either* Active, Inactive, or in Error; which of those states a component is in with respect to a particular context may be queried from the context itself.

#### 5.2.2.2.4 exit

#### Description

Stop the RTC's execution context(s) and finalize it along with its contents.

#### **Semantics**

Any execution contexts for which the RTC is the owner shall be stopped.

If the RTC participates in any execution contexts belonging to another RTC that contains it, directly or indirectly (i.e., the containing RTC is the **owner** of the **ExecutionContext**), it shall be deactivated in those contexts.

After the RTC is no longer Active in any Running execution context, it and any RTCs contained transitively within it shall be finalized.

#### **Constraints**

• An RTC cannot be exited if it has not yet been initialized. Any attempt to exit an RTC that is in the Created state shall fail with ReturnCode t::PRECONDITION NOT MET.

#### 5.2.2.2.5 attach context

#### **Description**

Inform this RTC that it is participating in the given execution context. Return a handle that represents the association of this RTC with the context.

#### **Semantics**

This operation is intended to be invoked by **ExecutionContextOperations::add\_component** (see Section 5.2.2.6.6). It is not intended for use by other clients.

#### 5.2.2.2.6 detach context

#### Description

Inform this RTC that it is no longer participating in the given execution context.

#### **Semantics**

This operation is intended to be invoked by **ExecutionContextOperations::remove\_component** (see Section 5.2.2.6.7). It is not intended for use by other clients.

#### **Constraints**

- This operation may not be invoked if this RTC is not already participating in the execution context. Such a call shall fail with ReturnCode t::PRECONDITION NOT MET.
- This operation may not be invoked if this RTC is Active in the indicated execution context. Otherwise, it shall fail with ReturnCode t::PRECONDITION NOT MET.

#### 5.2.2.2.7 get\_context

#### **Description**

Obtain a reference to the execution context represented by the given handle.

#### **Semantics**

The mapping from handle to context is specific to a particular RTC instance. The given handle must have been obtained by a previous call to **attach** context on this RTC.

#### 5.2.2.2.8 get\_owned\_contexts

### **Description**

This operation returns a list of all execution contexts owned by this RTC.

#### 5.2.2.2.9 get\_participating\_contexts

### Description

This operation returns a list of all execution contexts in which this RTC participates.

#### **Semantics**

Each call to attach\_context causes the provided context to be added to this list. Each call to detach\_context causes the provided context to be removed from this list.

## 5.2.2.2.10 get\_context\_handle

### Description

This operation returns a handle that is associated with the given execution context.

#### **Semantics**

The handle returned by this operation is same as the handle returned by attach\_context.

## 5.2.2.3 LifeCycleState

## Description

**LifeCycleState** is an enumeration of the states in the lifecycle above.

#### **Attributes**

LifeCycleState		
attributes		
CREATED	LifeCycleState	
INACTIVE	LifeCycleState	
ACTIVE	LifeCycleState	
ERROR	LifeCycleState	
no operations		

#### 5.2.2.3.1 CREATED

#### **Description**

The RTC object has been instantiated but not yet fully initialized.

## **5.2.2.3.2 INACTIVE**

#### Description

The RTC is Alive but is not being invoked in any execution context (see Section 5.2.2.5), regardless of whether the context is Running or not.

#### **Semantics**

An instance of this state exists for each execution context in which the RTC participates. If the RTC does not participate in any execution context, a single instance of this state exists.

#### 5.2.2.3.3 ACTIVE

#### **Description**

The RTC is Alive and will be invoked in the execution context if the context is Running.

#### **Semantics**

An instance of this state exists for each execution context in which the RTC participates. If the RTC does not participate in any execution context, this state shall never be observed.

### 5.2.2.3.4 ERROR

### **Description**

The RTC has encountered a problem in a given execution context and cannot continue functioning in that context without being reset.

## 5.2.2.4 ComponentAction

## **Description**

The ComponentAction interface provides callbacks corresponding to the execution of the lifecycle operations of LightweightRTObject (see Section 5.2.2.2) and ExecutionContext (see Section 5.2.2.5). An RTC developer may implement these callback operations in order to execute application-specific logic pointing response to those transitions.

#### **Semantics**

Clients of an RTC are not expected to invoke these operations directly; they are provided for the benefit of the RTC middleware implementation.

## **Operations**

ComponentAction				
no attributes				
operations				
on_initialize		ReturnCode_t		
on_finalize		ReturnCode_t		
on_startup		ReturnCode_t		
	exec_handle	ExecutionContextHandle_t		
on_shutdown		ReturnCode_t		
	exec_handle	ExecutionContextHandle_t		
on_activated		ReturnCode_t		
	exec_handle	ExecutionContextHandle_t		
on_deactivated		ReturnCode_t		
	exec_handle	ExecutionContextHandle_t		
on_aborting		ReturnCode_t		
	exec_handle	ExecutionContextHandle_t		
on_error		ReturnCode_t		

	exec_handle	ExecutionContextHandle_t
on_reset		ReturnCode_t
	exec_handle	ExecutionContextHandle_t

## 5.2.2.4.1 on\_initialize

## Description

The RTC has been initialized and entered the Alive state.

#### **Semantics**

Any RTC-specific initialization logic should be performed here.

## 5.2.2.4.2 on\_finalize

## Description

The RTC is being destroyed.

#### **Semantics**

Any final RTC-specific tear-down logic should be performed here.

## 5.2.2.4.3 on\_startup

## Description

The given execution context, in which the RTC is participating, has transitioned from Stopped to Running.

## 5.2.2.4.4 on\_shutdown

### Description

The given execution context, in which the RTC is participating, has transitioned from Running to Stopped.

## 5.2.2.4.5 on\_activated

### Description

The RTC has been activated in the given execution context.

## 5.2.2.4.6 on\_deactivated

## Description

The RTC has been deactivated in the given execution context.

## 5.2.2.4.7 on\_aborting

#### Description

The RTC is transitioning from the Active state to the Error state in some execution context.

#### **Semantics**

This callback is invoked only a single time for time that the RTC transitions into the Error state from another state. This behavior is in contrast to that of **on error**.

#### 5.2.2.4.8 on error

#### Description

The RTC remains in the Error state.

#### **Semantics**

If the RTC is in the Error state relative to some execution context when it would otherwise be invoked from that context (according to the context's **ExecutionKind**), this callback shall be invoked instead. For example,

- If the ExecutionKind is PERIODIC, this operation shall be invoked in sorted order at the rate of the context instead of DataFlowComponentAction::on execute and on state update.
- If the ExecutionKind is EVENT\_DRIVEN, this operation shall be invoked whenever FsmParticipantAction::on action would otherwise have been invoked.

## 5.2.2.4.9 on\_reset

#### Description

The RTC is in the Error state. An attempt is being made to recover it such that it can return to the Inactive state.

#### **Semantics**

If the RTC was successfully recovered and can safely return to the Inactive state, this method shall complete with ReturnCode\_t::OK. Any other result shall indicate that the RTC should remain in the Error state.

#### 5.2.2.5 ExecutionContext

#### Description

An **ExecutionContext** allows the business logic of an RTC to be decoupled from the thread of control in which it is executed. The context represents a logical thread of control and is provided to RTCs at runtime as an argument to various operations, allowing them to query and modify their own state, and that of other RTCs executing within the same context, in the lifecycle.

This separation of concerns is important for two primary reasons:

• Large number of components may collaborate tightly within a single node or process. If each component were to run within its own thread of control, the infrastructure may not be able to satisfy the timeliness and determinism requirements of real-time applications due to the large number of threads and the required synchronization between them.

• A single application may carry out a number of independent tasks that require different execution rates. For example, it may need to sample a sensor periodically at a very high rate and update a user interface at a much lower rate.

#### Interface Realizations

ExecutionContextOperations

#### **Semantics**

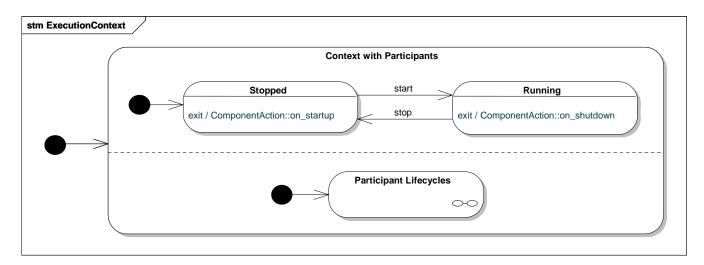


Figure 5.7 - ExecutionContext

The state machine of an **ExecutionContext** has two parts. The behavior of the **ExecutionContext** itself is defined by the upper region in the above figure. The behavior of the RTCs that participate in the context is defined by the lower region. The contents of that region are displayed in more detail in Figure 5.5 in Section 5.2.2.2.

#### Ownership and Participation

Each execution context is owned by a single RTC and may be used to execute that RTC and the RTCs contained within it, directly or indirectly. An RTC that owns one or more execution contexts is known as an *autonomous* RTC.

An autonomous RTC and some subset of the RTCs within it (to be defined by the application developer) shall be executed by the infrastructure according to the context's *execution kind*, which defines when each RTC's operations will be invoked when and in which order. These RTCs are said to *participate* in the context. The available execution kinds are described in Section 5.2.2.7.

The relationship between RTCs and execution contexts may be many-to-many in the general case: multiple RTCs may be invoked from the same execution context, and a single RTC may be invoked from multiple contexts. In the case where multiple RTCs are invoked from the same context, starting or stopping the context shall result in the corresponding lifecycle transitions for *all* of those components.

## Logical and Physical Threads

Although an execution context represents a logical thread of control, the choice of how it maps to a physical thread shall be left to the application's deployment environment. Implementations may elect to associate contexts with threads with a one-to-one mapping, to serve multiple contexts from a single thread, or by any other means. In the case where a given RTC may be invoked from multiple contexts, concurrency management is implementation-dependent.

## 5.2.2.6 ExecutionContextOperations

## **Description**

The ExecutionContextOperations interface defines the operations that all instances of ExecutionContext must provide.

## **Operations**

ExecutionContextOperations				
no attributes				
operations				
is_running		Boolean		
start		ReturnCode_t		
stop		ReturnCode_t		
get_rate		Double		
set_rate		ReturnCode_t		
	rate	Double		
add_component		ReturnCode_t		
	component	LightweightRTObject		
remove_component		ReturnCode_t		
	component	LightweightRTObject		
activate_component		ReturnCode_t		
	component	LightweightRTObject		
deactivate_component		ReturnCode_t		
	component	LightweightRTObject		
reset_component		ReturnCode_t		
	component	LightweightRTObject		
get_component_state		LifeCycleState		

	component	LightweightRTObject
get_kind		ExecutionKind

## 5.2.2.6.1 is\_running

## Description

This operation shall return true if the context is in the Running state.

#### **Semantics**

While the context is Running, all Active RTCs participating in the context shall be executed according to the context's execution kind.

## 5.2.2.6.2 start

### Description

Request that the context enter the Running state. Once the state transition occurs, the ComponentAction::on\_startup operation (see Section 5.2.2.4.3) will be invoked.

#### **Semantics**

An execution context may not be started until the RT components that participate in it have been initialized.

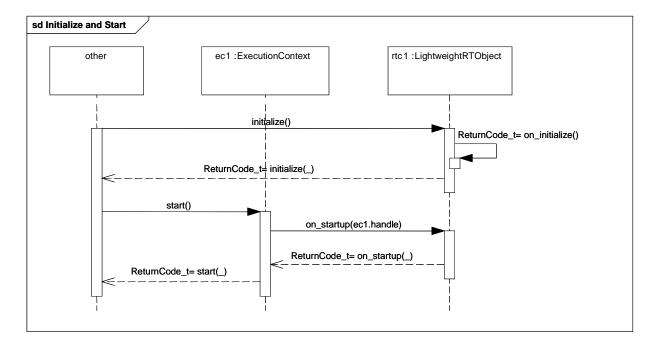


Figure 5.8 - Initialize and Start

An execution context may be started and stopped multiple times.

#### **Constraints**

- This operation shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET if the context is not in the Stopped state.
- This operation shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET if any of the participating components are not in their Alive state.

## 5.2.2.6.3 stop

### **Description**

Request that the context enter the Stopped state. Once the transition occurs, the ComponentAction::on\_shutdown operation (see Section 5.2.2.4.4) will be invoked.

#### **Semantics**

An execution context must be stopped before the RT components that participate in it are finalized.

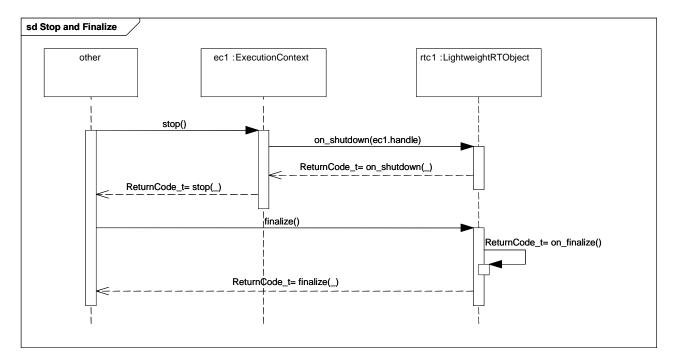


Figure 5.9 - Stop and Finalize

An execution context may be started and stopped multiple times.

### **Constraints**

• This operation shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET if the context is not in the Running state.

## 5.2.2.6.4 get\_rate

### **Description**

This operation shall return the rate (in hertz) at which its Active participating RTCs are being invoked.

#### **Semantics**

An implementation is permitted to perform some periodic or quasi-periodic processing within an execution context with an **ExecutionKind** other than **PERIODIC**. In such a case, the result of this operation is implementation-defined. If no periodic processing of any kind is taking place within the context, this operation shall fail as described in Section 5.2.1.

#### **Constraints**

• If the context has an **ExecutionKind** of **PERIODIC**, this operation shall return a rate greater than zero.

#### 5.2.2.6.5 set rate

#### Description

This operation shall set the rate (in hertz) at which this context's Active participating RTCs are being called.

#### **Semantics**

If the execution kind of the context is **PERIODIC**, a rate change shall result in the invocation of **on\_rate\_changed** on any RTCs realizing **DataFlowComponentAction** that are registered with any RTCs participating in the context.

An implementation is permitted to perform some periodic or quasi-periodic processing within an execution context with an **ExecutionKind** other than **PERIODIC**. If such is the case, and the implementation reports a rate from **get\_rate**, this operation shall set that rate successfully provided that the given rate is valid. If no periodic processing of any kind is taking place within the context, this operation shall fail with **ReturnCode\_t::UNSUPPORTED**.

#### **Constraints**

• The given rate must be greater than zero. Otherwise, this operation shall fail with ReturnCode\_t::BAD\_PARAMETER.

## 5.2.2.6.6 add\_component

#### **Description**

The operation causes the given RTC to begin participating in the execution context.

#### **Semantics**

The newly added RTC will receive a call to **LightweightRTComponent::attach\_context** (see Section 5.2.2.2.5) and then enter the Inactive state.

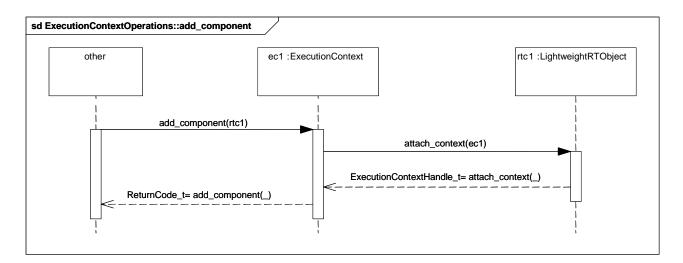


Figure 5.10 - ExecutionContextOperations::add\_component

- If the ExecutionKind of this context is **PERIODIC**, the RTC must be a data flow component (see Section 5.3.1.1). Otherwise, this operation shall fail with **ReturnCode** t::PRECONDITION NOT MET.
- If the ExecutionKind of this context is **EVENT\_DRIVEN**, the RTC must be an FSM participant (see Section 5.3.2.3). Otherwise, this operation shall fail with **ReturnCode t::PRECONDITION NOT MET**.

## 5.2.2.6.7 remove\_component

## **Description**

This operation causes a participant RTC to stop participating in the execution context.

### **Semantics**

The removed RTC will receive a call to LightweightRTComponent::detach context (see Section 5.2.2.2.6).

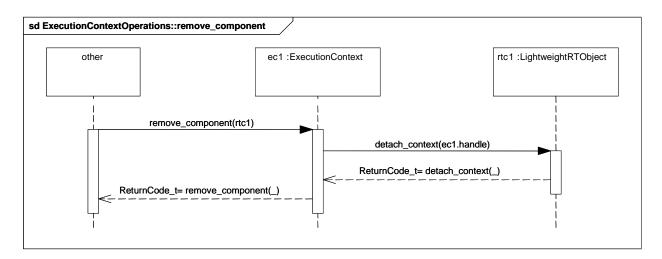


Figure 5.11 - ExecutionContextOperations::remove\_component

- If the given RTC is not currently participating in the execution context, this operation shall fail with ReturnCode\_t::BAD\_PARAMETER.
- An RTC must be deactivated before it can be removed from an execution context. If the given RTC is participating in the execution context but is still in the Active state, this operation shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET.

## 5.2.2.6.8 activate\_component

## Description

The given participant RTC is Inactive and is therefore not being invoked according to the execution context's execution kind. This operation shall cause the RTC to transition to the Active state such that it may subsequently be invoked in this execution context.

#### **Semantics**

The callback on\_activate shall be called as a result of calling this operation. This operation shall not return until the callback has returned, and shall result in an error if the callback does.

The following figure is a non-normative example sequence diagram for activate component.

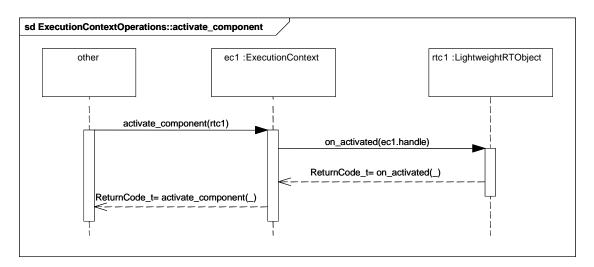


Figure 5.12 - ExecutionContextOperations::activate\_component

- An execution context can only activate its participant components. If the given RTC is not participating in the execution context, this operation shall fail with ReturnCode t::BAD PARAMETER.
- An RTC that is in the Error state cannot be activated until after it has been reset. If the given RTC is in the Error state, this operation shall fail with ReturnCode t::PRECONDITION NOT MET.
- This operation shall fail with ReturnCode t::BAD PARAMETER if the given component is not in its Alive state.

#### 5.2.2.6.9 deactivate\_component

## **Description**

The given RTC is Active in the execution context. Cause it to transition to the Inactive state such that it will not be subsequently invoked from the context unless and until it is activated again.

#### **Semantics**

The callback on\_deactivate shall be called as a result of calling this operation. This operation shall not return until the callback has returned, and shall result in an error if the callback does.

The following figure is a non-normative example sequence diagram for deactivate component.

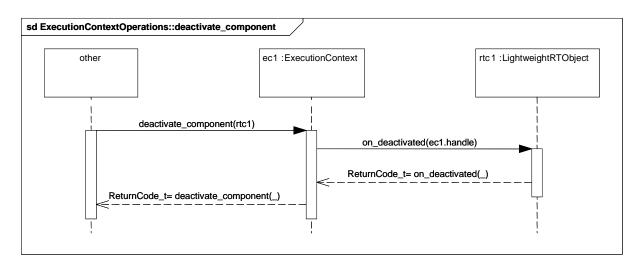


Figure 5.13 - ExecutionContextOperations::deactivate\_component

- An execution context can only deactivate its participant components. If the given RTC is not participating in the execution context, this operation shall fail with ReturnCode\_t::BAD\_PARAMETER.
- This operation shall fail with ReturnCode\_t::BAD\_PARAMETER if the given component is not in its Alive state.

#### 5.2.2.6.10 reset component

## Description

Attempt to recover the RTC when it is in Error.

## **Semantics**

The ComponentAction::on\_reset callback shall be invoked. This operation shall not return until the callback has returned, and shall result in an error if the callback does. If possible, the RTC developer should implement that callback such that the RTC may be returned to a valid state.

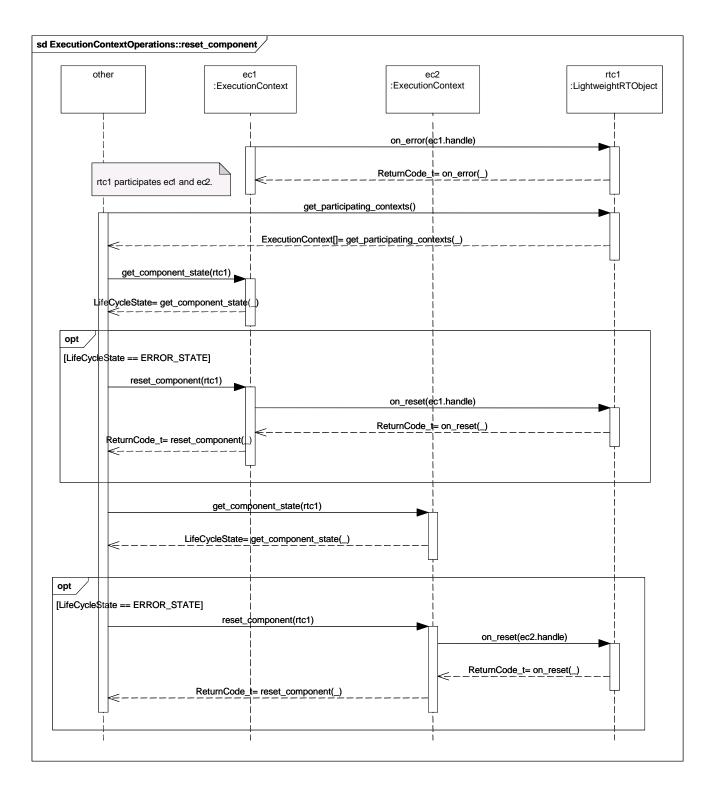


Figure 5.14 - ExecutionContextOperations::reset\_component

If this operation fails, the RTC will remain in Error.

#### **Constraints**

- An RTC may only be reset in an execution context in which it is in error. If the RTC is not in Error in the identified context, this operation shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET. However, that failure shall not cause the RTC to enter the Error state.
- An RTC may not be reset while in the Created state. Any attempt to invoke this operation while the RTC is in that state shall fail with ReturnCode\_t::PRECONDITION\_NOT\_MET. However, that failure shall not cause the RTC to enter the Error state.

## 5.2.2.6.11 get\_component\_state

### **Description**

This operation shall report the LifeCycleState of the given participant RTC.

#### **Constraints**

- The given RTC must be Alive.
- The given RTC must be a participant in the target execution context.
- The LifeCycleState returned by this operation shall be one of LifeCycleState::INACTIVE, ACTIVE, or ERROR.

### 5.2.2.6.12 get\_kind

#### Description

This operation shall report the execution kind of the execution context.

## 5.2.2.7 ExecutionKind

#### Description

The **ExecutionKind** enumeration defines the execution semantics (see Section 5.3) of the RTCs that participate in an execution context.

#### **Attributes**

ExecutionKind		
attributes		
PERIODIC	ExecutionKind	
EVENT_DRIVEN	ExecutionKind	
OTHER	ExecutionKind	
no operations		

## 5.2.2.7.1 PERIODIC

## **Description**

The participant RTCs are executing according to periodic sampled data semantics (see Section 5.3.1).

## 5.2.2.7.2 EVENT\_DRIVEN

## **Description**

The participant RTCs are executing according to stimulus response semantics (see Section 5.3.2).

## 5.2.2.7.3 OTHER

## **Description**

The participant RTCs are executing according to some semantics not defined by this specification.

## 5.2.2.8 ExecutionContextHandle\_t

## **Description**

This data type represents the association between an RTC and an ExecutionContext in which it participates.

#### **Semantics**

This is an opaque DataType. It has no attributes or operations.

# 5.2.3 Basic Types

This specification reuses the primitive types from [UML]: String, Boolean, etc. It also defines additional primitive types that are not available from UML, such as types for representing floating point numbers<sup>2</sup>. These additional types are described in this section.

#### 5.2.3.1 Character

### Description

The Character primitive type is an 8-bit quantity that encodes a single-byte character from any byte-oriented code set.

Character is an instance of PrimitiveType [UML].

#### **Constraints**

The Character value is a LiteralCharacter [UML].

#### 5.2.3.2 Double

#### **Description**

The **Double** primitive type is an IEEE double-precision floating point number. See IEEE Standard for Binary Floating-Point Arithmetic, ANSI/IEEE Standard 754-1985, for a detailed specification.

Double is an instance of PrimitiveType [UML].

#### **Constraints**

The Double value is a LiteralDouble.

#### 5.2.3.3 Float

#### **Description**

The **Float** primitive type is an IEEE single-precision floating point number. See IEEE Standard for Binary Floating-Point Arithmetic, ANSI/IEEE Standard 754-1985, for a detailed specification.

Float is an instance of PrimitiveType [UML].

#### **Constraints**

The Float value shall be a LiteralFloat that is an IEEE single-precision floating point number.

<sup>2.</sup> Note on model reuse: Models of the IDL basic types are also defined in the UML Profile for CORBA Specification [CORBAUML] and in the PIM and PSM for Software Radio Components document [SWRADIO] (although those two definitions differ from one another). CORBA and Software Radio Components are considered to be platforms by this specification; it is not appropriate for this PIM to depend on either of those specifications. Therefore, this specification defines the necessary types itself.

## 5.2.3.4 Long

#### Description

The **Long** primitive type, a specialization of **Integer** primitive type, is a signed integer range  $[-2^{31}, 2^{31} - 1]$ .

Long is an instance of PrimitiveType [UML].

#### **Constraints**

The Long value shall be a LiteralInteger [UML] with a range of  $[-2^{31}, 2^{31} - 1]$ .

## 5.2.3.5 LongDouble

### **Description**

The **LongDouble** primitive type is an IEEE double-extended floating-point number. See IEEE Standard for Binary Floating-Point Arithmetic, ANSI/IEEE Standard 754-1985, for a detailed specification.

LongDouble is an instance of PrimitiveType [UML].

#### **Constraints**

The LongDouble value shall be a LiteralLongDouble that is an IEEE double-extended floating-point number.

## 5.2.3.6 LongLong

## **Description**

The **LongLong** primitive type, a specialization of **Integer** primitive type, is a signed integer range  $[-2^{63}, 2^{63} - 1]$ . **LongLong** is an instance of **PrimitiveType** [UML].

### **Constraints**

The LongLong value shall be a LiteralInteger [UML] with a range of  $[-2^{63}, 2^{63} - 1]$ .

### 5.2.3.7 Octet

#### Description

The Octet primitive type, a specialization of Integer primitive type, is an unsigned integer within range [0, 255]. Octet is an instance of PrimitiveType [UML].

#### **Constraints**

The Octet value shall be a LiteralInteger [UML] with a range of [0, 255].

#### 5.2.3.8 Short

#### Description

The **Short** primitive type, a specialization of **Integer** primitive type, is a signed integer range  $[-2^{15}, 2^{15} - 1]$ .

Short is an instance of PrimitiveType [UML].

#### **Constraints**

The Short value shall be a LiteralInteger [UML] with a range of  $[-2^{15}, 2^{15} - 1]$ .

## 5.2.3.9 UnsignedLong

### **Description**

The **UnsignedLong** primitive type, a specialization of **Integer** primitive type, is an unsigned integer range  $[0, 2^{32} - 1]$ .

UnsignedLong is an instance of PrimitiveType [UML].

#### **Constraints**

The UnsignedLong value shall be a LiteralInteger [UML] with a range of  $[0, 2^{32} - 1]$ .

## 5.2.3.10 UnsignedLongLong

### Description

The **UnsignedLongLong** primitive type, a specialization of **Integer** primitive type, is an unsigned integer range  $[0, 2^{64} - 1]$ .

UnsignedLongLong is an instance of PrimitiveType [UML].

#### **Constraints**

The UnsignedLongLong value shall be a LiteralInteger [UML] with a range of [0, 2<sup>64</sup> - 1].

## 5.2.3.11 UnsignedShort

#### **Description**

The **UnsignedShort** primitive type, a specialization of **Integer** primitive type, is an unsigned integer range  $[0, 2^{16} - 1]$ .

UnsignedShort is an instance of PrimitiveType [UML].

### **Constraints**

The UnsignedShort value shall be a LiteralInteger [UML] with a range of [0, 2<sup>16</sup> - 1].

#### 5.2.3.12 WideCharacter

### Description

The WideCharacter primitive type represents a wide character that can be used for any character set.

WideCharacter is an instance of PrimitiveType [UML].

The WideCharacter value shall be a LiteralWideCharacter.

## 5.2.3.13 WideString

### **Description**

The WideString primitive type represents a wide character string that can be used for any character set.

WideString is an instance of PrimitiveType [UML].

#### **Constraints**

The WideString value is a LiteralWideString.

## 5.2.4 Literal Specifications

From [UML]:

A literal specification is an abstract specialization of **ValueSpecification** that identifies a literal constant being modeled.

#### 5.2.4.1 LiteralCharacter

## **Description**

A literal character contains a Character-valued attribute.

## Generalizations

• LiteralSpecification [UML]

#### **Attributes**

· value : Character

#### **Semantics**

A LiteralCharacter specifies a constant Character value.

#### 5.2.4.2 LiteralDouble

#### Description

A literal double contains a **Double**-valued attribute.

#### Generalizations

• LiteralSpecification [UML]

#### **Attributes**

· value: Double

#### **Semantics**

A LiteralDouble specifies a constant Double value.

#### 5.2.4.3 LiteralFloat

## Description

A literal float contains a Float-valued attribute.

#### Generalizations

• LiteralSpecification [UML]

#### **Attributes**

· value : Float

#### **Semantics**

A LiteralFloat specifies a constant Float value.

## 5.2.4.4 LiteralLongDouble

#### Description

A literal long double contains a LongDouble-valued attribute.

#### Generalizations

• LiteralSpecification [UML]

#### **Attributes**

• value : LongDouble

#### **Semantics**

A LiteralLongDouble specifies a constant LongDouble value.

## 5.2.4.5 LiteralWideCharacter

## Description

A literal wide character contains a WideCharacter-valued attribute.

#### Generalizations

• LiteralSpecification[UML]

#### **Attributes**

• value : WideCharacter

#### **Semantics**

A LiteralWideCharacter specifies a constant wide character value.

## 5.2.4.6 LiteralWideString

#### Description

A literal wide string contains a **WideString**-valued attribute.

#### Generalizations

• LiteralSpecification [UML]

#### **Attributes**

· value: WideString

#### **Semantics**

A LiteralWString specifies a constant wide string value.

## 5.3 Execution Semantics

Applications in many domains may benefit from a component-oriented design. Complex distributed control applications—including robotics applications—tend to share certain fundamental design patterns. Chief among them are periodic sampled data processing (also known as *data flow*), stimulus response processing (also known as *asynchronous* or *discrete event processing*), and modes of operation. This specification provides direct support for these design patterns in the sections that follow.

The semantics described in this section are applied to RTCs by means of stereotypes. These stereotypes are orthogonal and may be combined as desired within a single RTC. For example, a single RTC may participate in both periodic and stimulus response processing, or may execute periodically in multiple modes.

# 5.3.1 Periodic Sampled Data Processing

Real-time applications are frequently implemented using a methodology known as *periodic sampled data processing* or sometimes simply *data flow*<sup>3</sup>. This section defines the support for that design pattern and corresponds to the **PERIODIC** execution kind (see Section 5.2.2.7.1). A periodic execution context executes data flow components periodically in a well-defined order relative to one another. A data flow component may also be a composite component containing other data flow components. In this way, sampled data processing may be decomposed to an arbitrary level of hierarchy.

<sup>3.</sup> This methodology is supported by a number of existing software tools for real-time algorithm and/or application development, including Simulink from The MathWorks, LabView from National Instruments, and Constellation from RTI.

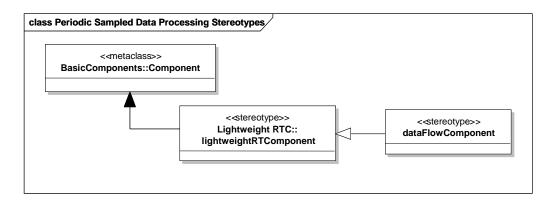


Figure 5.15 - Periodic Sampled Data Processing Stereotypes

An RTC developer declares a particular RTC to be a data flow component by extending it with the stereotype dataFlowComponent. The following figure is a non-normative example.

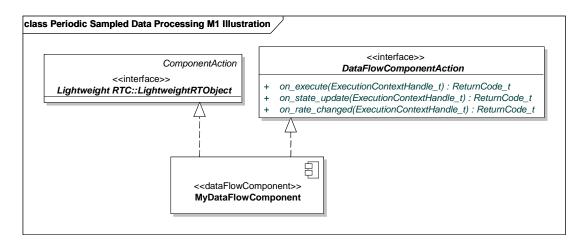


Figure 5.16 - Periodic Sampled Data Processing M1 Illustration

## 5.3.1.1 dataFlowComponent

#### **Description**

The dataFlowComponent stereotype may be applied to a component type to indicate that its instances should be executed in sorted order by a periodic execution context.

### **Constraints**

- An instance of a component extended by the dataFlowComponent stereotype must participate in at least one
  execution context of kind PERIODIC, which shall also be used for the execution of any contained data flow
  components.
- A component extended by dataFlowComponent must realize the interface DataFlowComponentAction.

#### Generalizations

• lightweightRTComponent

#### 5.3.1.1.1 Execution Sorting

A data flow component RTC determines the order in which to execute its contained data flow component RTCs by examining their interconnections in a process called *sorting*. The rules of sorting are simple: an RTC instance that produces an output (i.e., provides an interface) must run before any RTC instance that consumes that output (i.e., requires that interface). For example, if a provided interface of RTC instance A is connected to a required interface of RTC instance B, then A must be executed before B.

A data flow component RTC may contain component instances that are not data flow components. It shall ignore those instances for the purposes of periodic execution. However, other contained component instances may still invoke operations on the non-participant component instances if they choose.

Figure 5.17 depicts a non-normative example of a data flow component RTC with data flow components inside. The participants are executed in the order shown based upon their interconnections. The contained component that is not a data flow component is executed only when explicitly called by a data flow component.

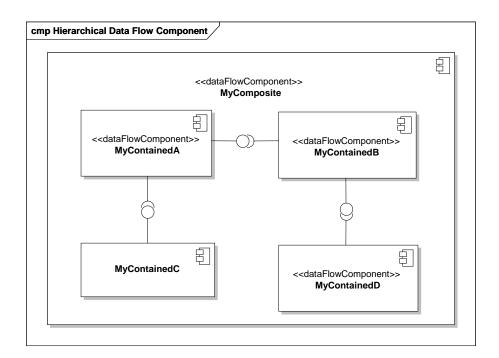


Figure 5.17 - Hierarchical Data flow Component

#### 5.3.1.1.2 Two-Pass Execution

In real-time applications, the need to minimize latency and jitter is critical. To reduce latency, the execution of each leaf RTC should occur as closely as possible to the beginning of the sample period. To reduce jitter, the execution of a particular leaf RTC should always occur at the same time relative to the beginning of the sample period. For these reasons, it is undesirable for an RTC to perform any task whose running time is long or indeterminate within on\_execute.

Furthermore, it is common for several collaborating RTCs to share access to a single data value or other resource. These RTCs should observe this resource in a consistent state during any sample period; i.e., it is inappropriate for any RTC to change the state of the resource before the other RTCs have executed.

In order to support the use cases just described (i.e., low latency, low jitter, and shared state), this specification supports two-pass execution of data flow component RTCs. In the first pass, each participant RTC's on\_execute (see Section 5.3.1.2.1) operation shall be invoked in sorted order. After the completion of the first pass, each RTC's on\_state\_update (see Section 5.3.1.2.2) operation shall be invoked in sorted order. For example, in Figure 5.17 the order of calls would be:

- 1. MyContainedA.on execute()
- 2. MyContainedB.on\_execute()
- 3. MyContainedD.on execute()
- 4. MyContainedA.on\_state\_update()
- 5. MyContainedB.on\_state\_update()
- 6. MyContainedD.on\_state\_update()

A data flow component should perform its primary business logic in **on\_execute** but delay any expensive operations or changes to shared state until **on state update**.

The following figures are non-normative example sequence diagrams for two-pass execution and ExecutionContextOperations::set rate.

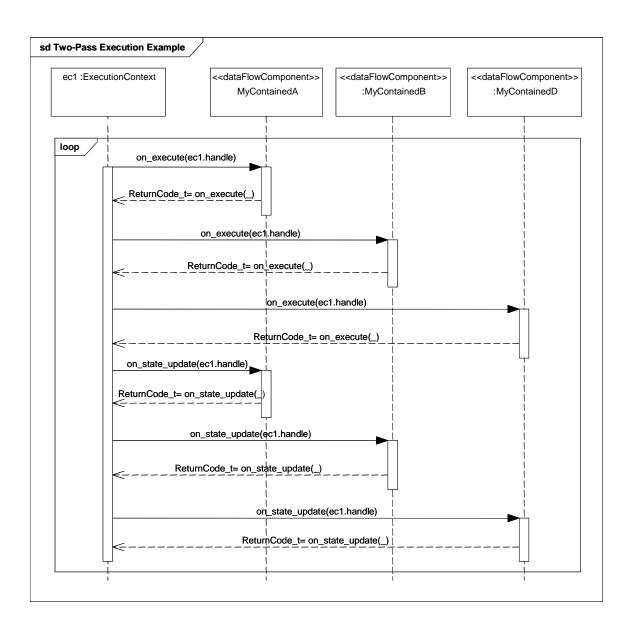


Figure 5.18 - Two Pass Execution Example

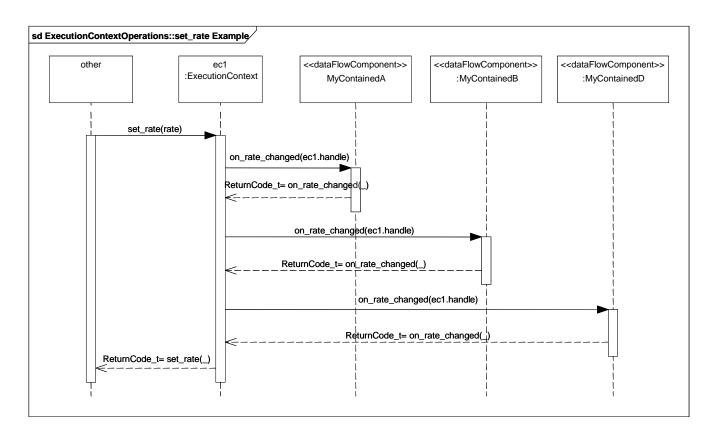


Figure 5.19 - ExecutionContextOperations::set\_rate Example

## 5.3.1.2 DataFlowComponentAction

## Description

**DataFlowComponentAction** is a companion to **ComponentAction** (see Section 5.2.2.4) that provides additional callbacks for intercepting the two execution passes defined in Section 5.3.1.1.2.

### **Operations**

DataFlowComponentAction		
no attributes		
operations		
on_execute		ReturnCode_t
	exec_handle	ExecutionContextHandle_t
on_state_update		ReturnCode_t
	exec_handle	ExecutionContextHandle_t

on_rate_changed		ReturnCode_t
	exec_handle	ExecutionContextHandle_t

#### 5.3.1.2.1 on\_execute

## **Description**

This operation will be invoked periodically at the rate of the given execution context as long as the following conditions hold:

- The RTC is Active.
- The given execution context is Running.

#### **Semantics**

This callback occurs during the first execution pass.

#### **Constraints**

• The execution context of the given context shall be **PERIODIC**.

## 5.3.1.2.2 on\_state\_update

### **Description**

This operation will be invoked periodically at the rate of the given execution context as long as the following conditions hold:

- The RTC is Active.
- The given execution context is Running.

### **Semantics**

This callback occurs during the second execution pass.

#### **Constraints**

• The execution context of the given context shall be **PERIODIC**.

## 5.3.1.2.3 on\_rate\_changed

## **Description**

This operation is a notification that the rate of the indicated execution context (see Section 5.2.2.6.4) has changed.

#### **Constraints**

• The execution context of the given context shall be **PERIODIC**.

## 5.3.2 Stimulus Response Processing

The *stimulus response* design pattern is also referred to as *asynchronous* or *discrete event processing*. Applications use this pattern when they need to respond asynchronously to changes in their environment. The behavior is modeled using finite state machines. The state machine waits until the asynchronous arrival of an "event", at which point it wakes up and transitions to a new state, executing an action associated with that transition.

This design pattern corresponds to the **EVENT DRIVEN** execution kind (see Section 5.2.2.7.2).

Stimulus response consists of two roles:

- A finite state machine composite component (FSM) contains a set of named states and transitions between them. Each
  transition, as well as the entry and exit of each state, may be associated with the execution of a contained finite state
  machine participant instance.
- A finite state machine participant declares to a containing FSM that it is available for execution when a transition is taken or when a state is entered or exited.

In many cases, the FSM will also be an FSM participant, allowing it to appear as a submachine state within another FSM.

Stimulus response models can be inspected and modified dynamically using types in the Introspection package. The correspondence between state machine transitions, state entries, and state exits in an FSM and FSM participants is described in an FsmProfile (see Section 5.4.1.7) and its FsmBehaviorProfiles (see Section 5.4.1.8) in the Introspection package (see Section 5.4). This data is provided by an FsmService interface (see Section 5.4.2.5) in the Introspection package.

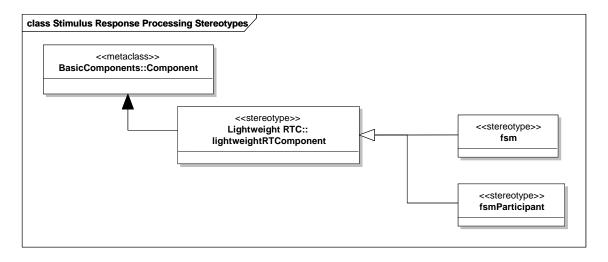


Figure 5.20 - Stimulus Response Processing Stereotypes

The RTC developer declares a particular RTC to be an FSM, or FSM participant, by extending it with the stereotype fsm (see Section 5.3.2.1) or fsmParticipant (see Section 5.3.2.3) respectively.

<sup>4.</sup> An Event in the nomenclature of finite state machines [UML] is not to be confused with "events" as understood within the context of the CORBA Event Service. Any relationship between the two is implementation-defined.

The following figure is a non-normative example of a UML model that incorporates these stereotypes.

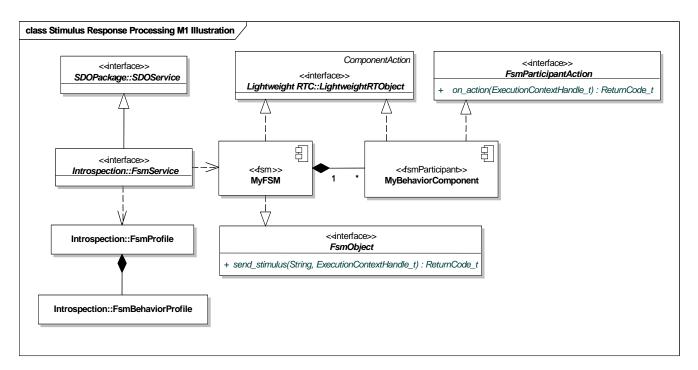


Figure 5.21 - Stimulus Response Processing M1 Illustration

## 5.3.2.1 fsm

## **Description**

Applying the fsm stereotype to a component implies the ability to define component-specific states and transitions.

## **Semantics**

In creating a state machine such as is depicted in Figure 5.22, the RTC developer is implicitly defining the Active state to be a submachine state.

The BehaviorStateMachines package described in [UML] is considered the normative definition of a state machine.

### Generalizations

• lightweightRTComponent

#### **Constraints**

- A component extended by fsm must realize the interface FsmObject.
- A component extended by the fsm stereotype must participate in at least one execution context of kind EVENT DRIVEN.
- The isSubmachineState attribute of the Active state must be true.

• The submachine association of the Active state must refer to a non-nil StateMachine.

## 5.3.2.2 FsmObject

#### Description

The FsmObject interface allows programs to send stimuli to a finite state machine, possibly causing it to change states.

#### **Operations**

${\it FsmObject}$		
no attributes		
operations		
send_stimulus		ReturnCode_t
	message	String
	exec_handle	ExecutionContextHandle_t

## 5.3.2.2.1 send\_stimulus

#### **Description**

Send a stimulus to an FSM that realizes this interface.

## **Semantics**

If the stimulus corresponds to any outgoing transition of the current state, that transition shall be taken and the state shall change. Any FSM participants associated with the exit of the current state, the transition to the new state, or the entry to the new state shall be invoked. If the stimulus does not correspond to any such transition, this operation shall succeed but have no effect.

If the given execution context is a non-nil reference to a context in which this FSM participates, the transition shall be executed in that context. If the argument is nil, the FSM shall choose an **EVENT\_DRIVEN** context in which to execute the transition. If the argument is non-nil, but this FSM does not participate in the given context, this operation shall fail with **ReturnCode\_t::BAD\_PARAMETER**.

#### **Constraints**

• The given execution context shall be of kind **EVENT DRIVEN**.

## 5.3.2.3 fsmParticipant

#### Description

Applying the **fsmParticipant** stereotype to a component implicitly defines a **Behavior** [UML] of the same name as the RTC instance itself that may be invoked by a containing FSM when transitioning between states, entering or exiting a state, and so forth. The invocation of the **Behavior** corresponds to the invocation of the **on\_action** (see Section 5.3.2.4.1) operation of the FSM participant RTC.

Figure 5.22 is a non-normative example of how a **Behavior** defined in terms of an FSM participant might be displayed graphically.

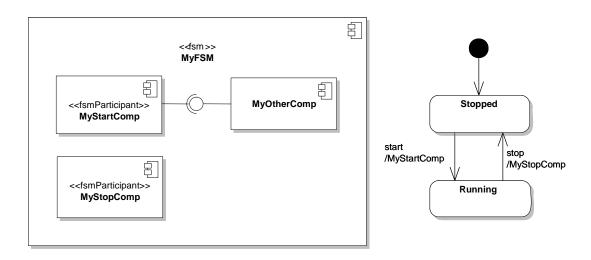


Figure 5.22 - FSM Participant Defines State Transition Behavior

## **Semantics**

The following figure is a non-normative example sequence diagram for Stimulus Response Processing.

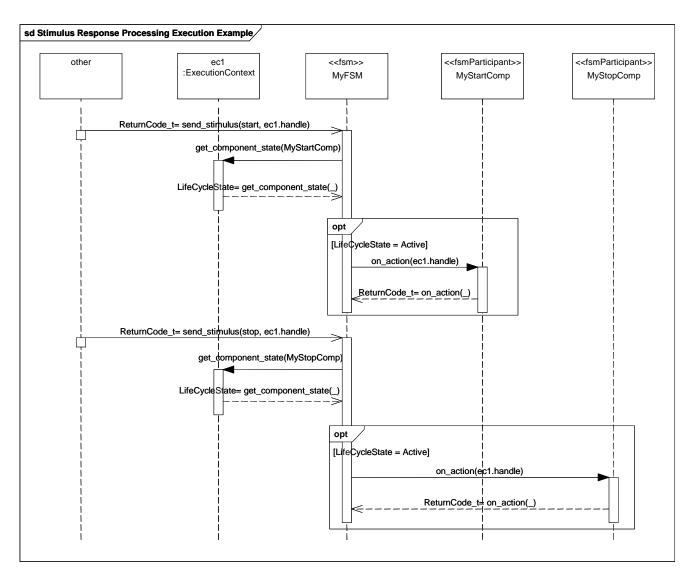


Figure 5.23 - Stimulus Responding Processing Execution Example

- A component extended by the fsmParticipant stereotype must realize the interface FsmParticipantAction.
- A component extended by the fsmParticipant stereotype must participate in at least one execution context of kind EVENT DRIVEN.

#### Generalizations

• lightweightRTComponent

## 5.3.2.4 FsmParticipantAction

### **Description**

**FsmParticipantAction** is companion to **ComponentAction** (see Section 5.2.2.4) that is intended for use with FSM participant RTCs. It adds a callback for the interception of state transitions, state entries, and state exits.

#### **Operations**

FsmParticipantAction		
no attributes		
operations		
on_action		ReturnCode_t
	exec_handle	ExecutionContextHandle_t

#### 5.3.2.4.1 on\_action

## **Description**

The indicated FSM participant RTC has been invoked as a result of a transition, state entry, or state exit in its containing FSM.

#### **Constraints**

• The given execution context shall be of kind EVENT DRIVEN.

# 5.3.3 Modes of Operation

Modes of operation provide support for applications that need to switch between different implementations of a given functionality. For example, an automobile may throttle its engine either based solely on the position of the gas pedal or alternatively based on the desired speed set by the cruise control depending on whether the cruise control is activated or not. "Cruise control on" and "cruise control off" are examples of modes. (In this example, the choice of mode is a binary one, although in general any number of alternatives may be necessary.)

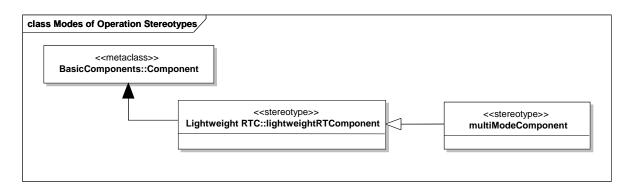


Figure 5.24 - Mode of Operation Stereotypes

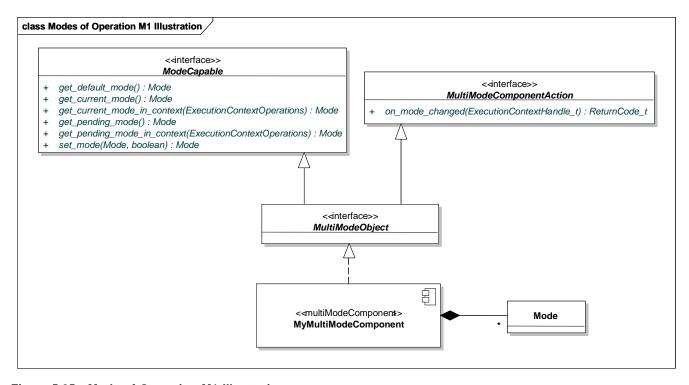


Figure 5.25 - Mode of Operation M1 Illustration

## 5.3.3.1 multiModeComponent

## Description

An RTC developer may declare an RTC to support multiple modes by extending it with the stereotype multiModeComponent, which introduces the operations defined in the interface MultiModeObject and its supertypes.

#### Generalizations

• lightweightRTComponent

#### **Constraints**

• A component extended by the multiModeComponent stereotype must realize the MultiModeObject interface.

## 5.3.3.2 MultiModeObject

### Description

The MultiModeObject interface and its supertypes define the operations that must be provided by any RTC that supports multiple modes.

#### Generalizations

- ModeCapable
- MultiModeComponentAction

## 5.3.3.3 ModeCapable

#### Description

The ModeCapable interface provides access to an object's modes and a means to set the current mode.

### **Semantics**

A given RTC may support multiple modes as well as multiple execution contexts. In such a case, a request for a mode change (e.g., from "cruise control on" to "cruise control off") may come asynchronously with respect to one or more of those execution contexts. The mode of an RTC may therefore be observed to be different from one execution context to another.

- A mode is pending in a given execution context when a mode change has been requested but the new mode has not yet been observed by that context.
- The new mode has been committed in a given execution context when the context finally observes the new mode.
- The new mode has stabilized once it has been committed in all execution contexts in which the RTC participates.

Figure 5.26 depicts a state machine that describes mode changes. Each parallel region in the composite state Mode Pending represents an execution context. The trigger "sample" within that state is considered to have occurred:

- ...just before the next call to on\_execute (see Section 5.3.1.2.1) in the case where immediate is false and the execution kind is PERIODIC, ...
- ...just before the processing of the next stimulus in the case where **immediate** is **false** and the execution kind is **EVENT DRIVEN**, or ...
- ...immediately in all other cases.

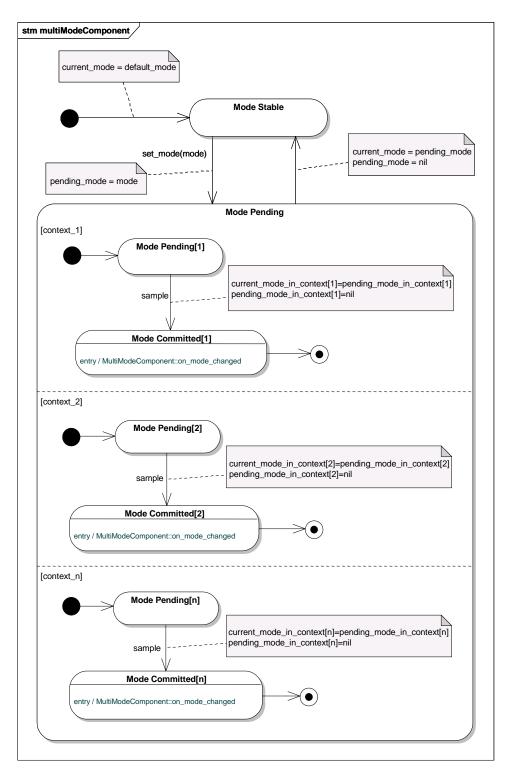


Figure 5.26 - multiModeComponent State Machine

The following figure shows a mode change on a component that is multi-mode and also a data flow participant. The mode change is not immediate, so it waits for the next "tick" of the execution context. This is a non-normative example.

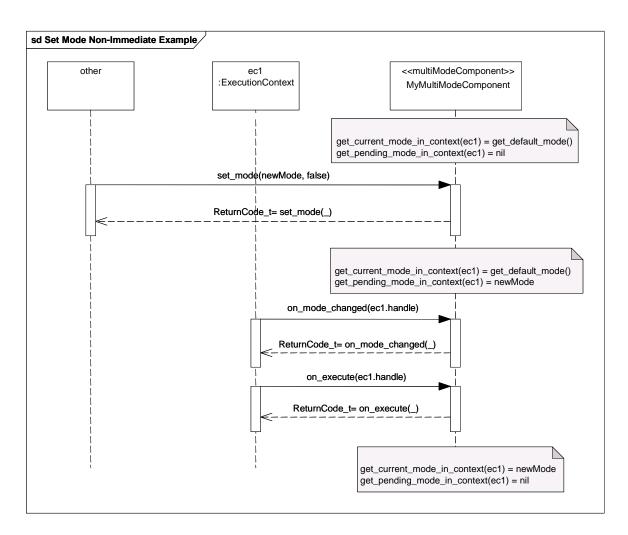


Figure 5.27 - Set Mode Non-Immediate Example

The following figure shows the same mode change, but this time it happens immediately instead of waiting for the execution context.

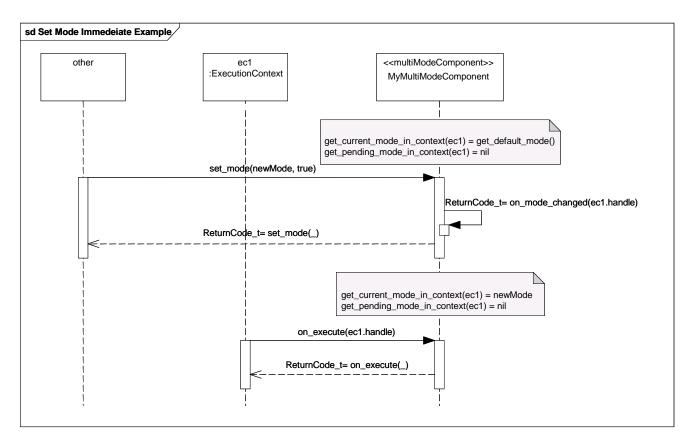


Figure 5.28 - Set Mode Immediate Example

## **Operations**

ModeCapable		
no attributes		
operations		
get_default_mode		Mode
get_current_mode		Mode
get_current_mode_in_context		Mode
	context	ExecutionContextOperations
get_pending_mode		Mode
get_pending_mode_in_context		Mode
	context	ExecutionContextOperations
set_mode		ReturnCode_t
	mode	Mode
	immediate	Boolean

## 5.3.3.3.1 get\_default\_mode

## **Description**

This operation shall return the mode in which the RTC shall be when no other mode has been set.

## **Constraints**

• This operation shall not return nil.

## 5.3.3.3.2 get\_current\_mode

## **Description**

This operation shall return the last mode to have stabilized. If no mode has been explicitly set, the current mode shall be the default mode.

## **Constraints**

• This operation shall never return nil.

## 5.3.3.3 get\_current\_mode\_in\_context

## **Description**

This operation returns the current mode of the component as seen by the indicated execution context.

#### **Semantics**

The manner in which this property changes is described in Figure 5.26.

#### 5.3.3.3.4 get\_pending\_mode

### Description

This operation shall return the last mode to have been passed to **set\_mode** that has not yet stabilized. Once the RTC's mode has stabilized, this operation shall return nil.

## 5.3.3.5 get\_pending\_mode\_in\_context

## Description

If the last mode to be requested by a call to **set\_mode** is different than the current mode as seen by the indicated execution context (see **get\_current\_mode\_in\_context**), this operation returns the former. If the requested mode has already been seen in that context, it returns nil.

#### **Semantics**

See Figure 5.26 for a description of how the pending mode relates to the current mode within a given execution context.

#### 5.3.3.3.6 set mode

### Description

This operation shall request that the RTC change to the indicated mode.

#### **Semantics**

Usually, the new mode will be *pending* in each execution context in which the component executes until the next sample period (if the execution kind is PERIODIC); at that point it will become the *current* mode in that context and there will no longer be a pending mode. However, in some cases it is important for a mode change to take place immediately; for example, a serious fault has occurred and the component must enter an emergency mode to ensure fail-safe behavior in a safety-critical system. In such a case, **immediate** should be **true** and the mode change will take place in all contexts without waiting for the next sample period.

### 5.3.3.4 MultiModeComponentAction

MultiModeComponentAction is a companion to ComponentAction that is realized by RTCs that support multiple modes.

${\it MultiModeComponentAction}$		
no attributes		
operations		
on_mode_changed		ReturnCode_t
	exec_handle	ExecutionContextHandle_t

## 5.3.3.4.1 on\_mode\_changed

### **Description**

This callback is invoked each time the observed mode of a component has changed with respect to a particular execution context.

#### **Semantics**

If the context is **PERIODIC**, this callback shall come before the next call to **on\_execute** (see Section 5.3.1.2.1) within that context.

The new mode can be retrieved with get\_current\_mode\_in\_context. If the result is the same as the result of get current mode, the mode has stabilized.

### 5.3.3.5 Mode

#### Description

Each mode defined by a given RTC shall be represented by an instance of Mode.

# 5.4 Introspection

This section of the PIM extends the Lightweight RTC model and Super Distributed Objects [SDO]. It describes a capability for querying and administering RTCs at runtime. These capabilities may be used by other RTCs—in implementations that support dynamic RTC composition—but also by tools or other supporting technologies.

Introspection functionalities and "lightweight" functionalities are described separately in this specification. This design has advantages, especially for applications targeting resource-constrained environments:

- The introspection interfaces may be remote (in implementations that support that capability) without requiring that every component and supporting object be available remotely.
- Lightweight RTCs can have a much lighter footprint and can dedicate themselves more completely to their business function rather than to ancillary functionality, introspection being just one example.
- This section is divided into two subsections:
  - Resource Data Model
  - Stereotype and Interfaces

The former specifies data-only classes called "profiles" that act as descriptors for the important entities described in this document. The latter describes the behavioral classes that provide those profiles.

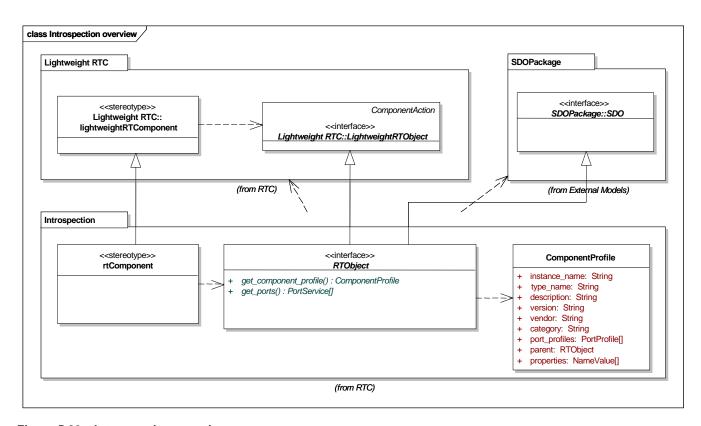


Figure 5.29 - Introspection overview

## 5.4.1 Resource Data Model

This section specifies the RTC resource data model, a set of data-only classes used to describe the capabilities and properties of RTCs. The RTC resource data model consists of data structure known as "profiles" that act as descriptors for the significant entities defined by the earlier portions of this specification. The data they contain may be "live" or may be populated using design-time information.

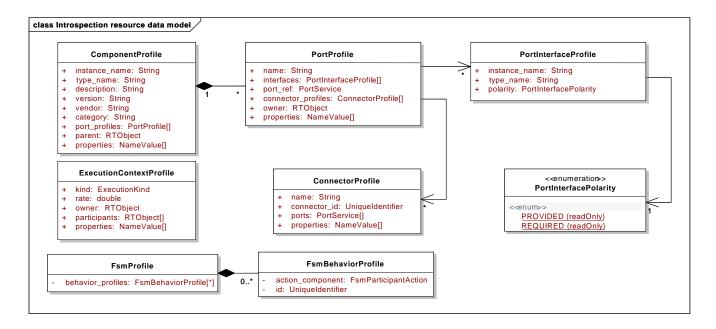


Figure 5.30 - Introspection Resource Data Model

## 5.4.1.1 ComponentProfile

## **Description**

ComponentProfile represents the static state of an RTC that is referred to here as the "target" RTC.

## **Attributes**

ComponentProfile		
attributes		
instance_name	String	
type_name	String	
description	String	
version	String	
vendor	String	
category	String	

port_p	rofiles	PortProfile[]
	parent	RTObject
pro	perties	NameValue[]
no operations		

## 5.4.1.1.1 instance\_name

## Description

This attribute shall contain the name of the target RTC instance.

#### **Semantics**

The instance name should be unique among RTC instances contained within the same containing component.

## 5.4.1.1.2 type\_name

## **Description**

This attribute shall contain the name of the target RTC class.

## **Semantics**

Each RTC class must have a name that is unique within an application.

## 5.4.1.1.3 description

#### **Description**

This attribute shall briefly describe the target RTC for the benefit of a human operator.

#### 5.4.1.1.4 version

## Description

This attribute shall contain the version number of the target RTC class.

#### **Semantics**

The format of the version number is outside of the scope of this specification.

## 5.4.1.1.5 vendor

## Description

The name of the individual or organization that produced the target RTC class.

## 5.4.1.1.6 category

#### **Description**

This attribute contains the name of a "category" or group to which the target RTC belongs.

## 5.4.1.1.7 port\_profiles

#### Description

This attribute contains a list of **PortProfiles** that describe the ports of the target RTC.

#### **Semantics**

There shall be a one-to-one correspondence between the members of this list and the ports of the target RTC.

## 5.4.1.1.8 parent

#### Description

This attribute contains a reference to the RTC that contains the target RTC instance. If the target RTC instance is not owned by any other RTC, this field stores a nil reference.

## **5.4.1.1.9 properties**

#### Description

This attribute contains additional properties of the target RTC.

#### **Semantics**

This attribute provides implementations the opportunity to describe additional characteristics of a particular RTC that are otherwise outside of the scope of this specification.

## 5.4.1.2 PortProfile

## **Description**

A PortProfile describes a port of an RTC (referred to as the "target" RTC). This port is referred to as the "target" port.

From this profile, other components and tools can obtain Port's name, type, object reference, and so on.

#### **Attributes**

PortProfile			
attributes			
name	String		
interfaces	PortInterfaceProfile[]		
port_ref	PortService		
connector_profiles	ConnectorProfile[]		
owner	RTObject		
properties	NameValue[]		
no operations			

## 5.4.1.2.1 name

## Description

This attribute contains the name of the target port.

## **Semantics**

Ports owned by an RTC are distinguished by their names. Therefore, this name should be unique within the target RTC.

## 5.4.1.2.2 interfaces

## Description

This attribute contains the name and polarity of each interface exposed by the target port.

## 5.4.1.2.3 port\_ref

## **Description**

This attributes contains a reference to the target port.

## 5.4.1.2.4 connector\_profiles

## Description

This attribute contains a collection of profiles describing the connections to the target port.

## 5.4.1.2.5 owner

## Description

This attribute contains a reference to the target RTC.

## 5.4.1.2.6 properties

#### Description

This attribute contains additional properties of the port.

#### **Semantics**

This attribute provides implementations the opportunity to describe additional characteristics of a particular port that are otherwise outside of the scope of this specification.

## 5.4.1.3 PortInterfaceProfile

## **Description**

**PortInterfaceProfile** describes an instance of a particular interface as it is exposed by a particular port. These objects are referred to below as the "target interface" and "target port" respectively.

#### **Attributes**

PortInterfaceProfile		
attributes		
instance_name	String	
type_name	String	
polarity	PortInterfacePolarity	
no operations		

## 5.4.1.3.1 instance\_name

## **Description**

This attribute stores the name of the target interface instance.

## 5.4.1.3.2 type\_name

## **Description**

This attribute stores the name of the target interface type.

## 5.4.1.3.3 polarity

## **Description**

This attribute indicates whether the target interface instance is provided or required by the RTC.

# 5.4.1.4 PortInterfacePolarity

## Description

The PortInterfacePolarity enumeration identifies exposed interface instances as provided or required.

#### **Attributes**

PortInterfacePolarity		
attributes		
PROVIDED	PortInterfacePolarity	
REQUIRED	PortInterfacePolarity	
no operations		

## 5.4.1.4.1 PROVIDED

## Description

The target interface is provided as an output by the target port.

## 5.4.1.4.2 REQUIRED

#### **Description**

The target interface is required as an input by the target port.

## 5.4.1.5 ConnectorProfile

## **Description**

The ConnectorProfile contains information about a connection between the ports of collaborating RTCs.

#### **Attributes**

ConnectorProfile			
attributes			
name	String		
connector_id	UniqueIdentifier		
ports	PortService[]		
properties	NameValue[]		
no operations			

## 5.4.1.5.1 name

## **Description**

This attribute contains the name of this connection.

## 5.4.1.5.2 connector\_id

## **Description**

Each connector has a unique identifier that is assigned when connection is established. This attribute stores that identifier.

## 5.4.1.5.3 ports

## **Description**

This field stores references to all ports connected by the target connector.

## **5.4.1.5.4** properties

#### Description

This attribute contains additional properties of the connection.

#### **Semantics**

This attribute provides implementations the opportunity to describe additional characteristics of a particular connection that are outside of the scope of this specification.

## 5.4.1.6 ExecutionContextProfile

#### **Attributes**

	ExecutionContextProfile
attributes	

kind	ExecutionKind
rate	Double
owner	RTObject
participants	RTObject[]
properties	NameValue[]
no operations	

#### 5.4.1.6.1 kind

## Description

This attribute stores the context's **ExecutionKind**.

#### 5.4.1.6.2 rate

## Description

This attribute stores execution rate.

#### **Semantics**

If the execution kind is not **PERIODIC**, the value here may not be valid (and should be negative in that case). See **ExecutionContext::get\_rate** (see Section 5.2.2.6.4) and **set\_rate** (see Section 5.2.2.6.5) for more information.

#### 5.4.1.6.3 owner

#### Description

This attribute stores a reference to the RTC that owns the context.

## 5.4.1.6.4 participants

## Description

This attribute stores references to the context's participant RTCs.

## **5.4.1.6.5** properties

## Description

This attribute contains additional properties of the execution context.

#### **Semantics**

This attribute provides implementations the opportunity to describe additional characteristics of a particular execution context that are outside the scope of this specification.

## 5.4.1.7 FsmProfile

#### Description

The FsmProfile describes the correspondence between an FSM and its contained FSM participants. This Profile is necessary for Stimulus Response Processing.

#### **Attributes**

FsmProfile				
attributes	uttributes			
	behavior_profiles FsmBehaviorProfile[]			
no operations				

## 5.4.1.7.1 behavior\_profiles

#### Description

This attribute lists the correspondences between an FSM and its contained FSM participants.

## 5.4.1.8 FsmBehaviorProfile

## **Description**

FsmBehaviorProfile represents the association of an FSM participant with a transition, state entry, or state exit in an FSM.

## **Semantics**

The assignment of identifiers to particular transitions, state entries, or state exits is implementation-dependent.

#### **Attributes**

FsmBehaviorProfile			
attributes			
action_component FsmParticipantAction			
id UniqueIdentifier			
no operations			

## 5.4.1.8.1 action\_component

## **Description**

This attribute stores a reference to the FSM participant that is invoked when the containing Fsm receives a message distinguished by id.

## 5.4.1.8.2 id

## Description

This attribute stores the message identifier.

# 5.4.2 Stereotypes and Interfaces

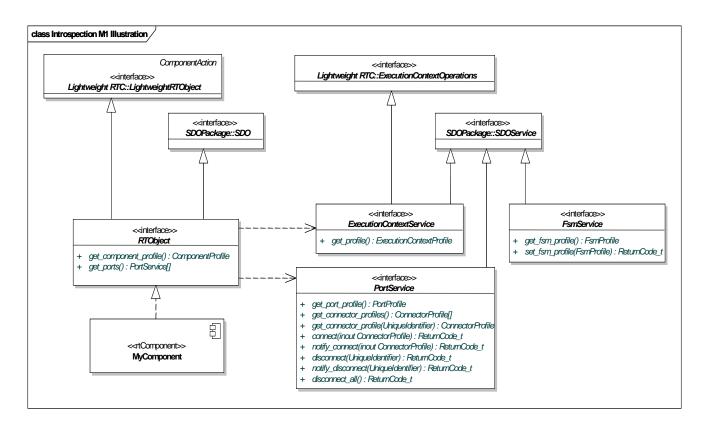


Figure 5.31 - Introspection M1 Illustration

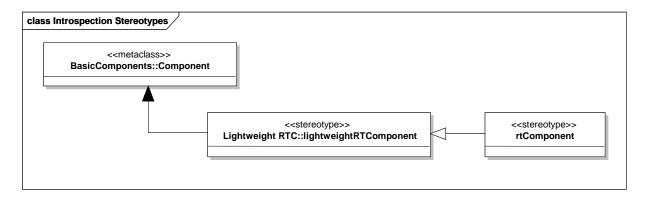


Figure 5.32 - Introspection Stereotypes

#### 5.4.2.1 rtComponent

## **Description**

The **rtComponent** stereotype identifies the component it extends as an RTC that realizes an SDO-based introspective interface.

#### Generalizations

• lightweightRTComponent

#### **Constraints**

• A component extended by the rtComponent stereotype must realize the RTObject interface.

## 5.4.2.2 RTObject

## **Description**

The RTObject interface defines the operations that all SDO-based RTCs must provide. It is required by the rtComponent stereotype.

## Generalizations

- LightweightRTObject
- SDO [SDO]

#### **Constraints**

Any execution contexts returned from the inherited methods
 LightweightRTObject::get\_owned\_contexts and
 LightweightRTObject::get\_participating\_contexts shall be of type
 ExecutionContextService.

## **Operations**

RTObject			
no attributes			
operations			
get_component_profile			ComponentProfile
get_ports			PortService[]

#### 5.4.2.2.1 get\_component\_profile

#### Description

This operation returns the ComponentProfile of the RTC.

## 5.4.2.2.2 get\_ports

## Description

This operation returns a list of the RTCs ports.

## 5.4.2.3 PortService

#### Description

An instance of the PortService interface represents a port (i.e., UML::Composite Structures::Ports::Port) of an RTC. It provides operations that allow it to be connected to and disconnected from other ports.

#### **Semantics**

A port service can support unidirectional or bidirectional communication.

A port service may allow for a service-oriented connection, in which other connected ports, invoke methods on it. It may also allow for a data-centric connection, in which data values are streamed in or out. In either case, the connection is described by an instance of **ConnectorProfile**. However, the behavioral contracts of such connections are dependent on the interfaces exposed by the ports and are not described normatively by this specification.

#### Generalizations

• SDOService [SDO]

## **Operations**

PortService PortService		
no attributes		
operations		

get_port_profile		PortProfile
get_connector_profiles		ConnectorProfile[]
get_connector_profile		ConnectorProfile
	connector_id	UniqueIdentifier
connect		ReturnCode_t
	connector_profile	ConnectorProfile
notify_connect		ReturnCode_t
	inout connector_profile	ConnectorProfile
disconnect		ReturnCode_t
	connector_id	UniqueIdentifier
notify_disconnect		ReturnCode_t
	connector_id	UniqueIdentifier
disconnect_all		ReturnCode_t

## 5.4.2.3.1 get\_port\_profile

## **Description**

This operation returns the **PortProfile** of the **PortService**.

## 5.4.2.3.2 get\_connector\_profiles

#### **Description**

This operation returns a list of the ConnectorProfiles of the PortService.

## 5.4.2.3.3 get\_connector\_profile

#### **Description**

This operation returns the **ConnectorProfile** specified by a connector ID.

#### 5.4.2.3.4 connect

## **Description**

This operation establishes connection between this port and the peer ports according to given ConnectionProfile.

#### **Semantics**

A ConnectorProfile has a sequence of port references. This port invokes the notify\_connect operation of one of the ports included in the sequence. It follows that the notification of connection is propagated by the notify\_connect operation with ConnectorProfile. This operation returns ConnectorProfile return value and returns ReturnCode t as return codes.

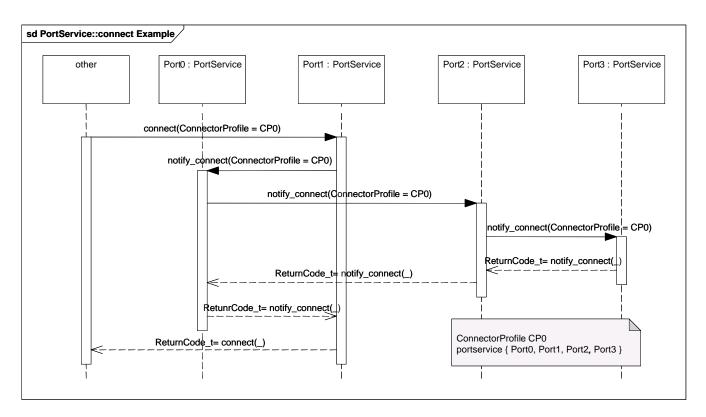


Figure 5.33 - PortService::connect Example

#### 5.4.2.3.5 disconnect

## **Description**

This operation destroys the connection between this port and its peer ports using the ID that was given when the connection was established.

#### **Semantics**

This port invokes the **notify\_disconnect** operation of one of the ports included in the sequence of the **ConnectorProfile** stored when the connection was established. The notification of disconnection is propagated by the **notify disconnect** operation.

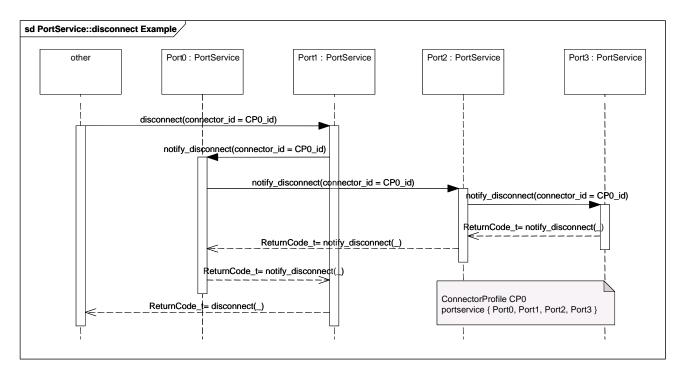


Figure 5.34 - PortService::disconnect Example

#### 5.4.2.3.6 disconnect all

#### Description

This operation destroys all connection channels owned by the PortService.

## 5.4.2.3.7 notify\_connect

## **Description**

This operation notifies this **PortService** of the connection between its corresponding port and the other ports and propagates the given **ConnectionProfile**.

#### **Semantics**

A ConnectorProfile has a sequence of port references. This PortService stores the ConnectorProfile and invokes the notify\_connect operation of the next PortService in the sequence. As ports are added to the connector, PortService references are added to the ConnectorProfile and provided to the caller. In this way, notification of connection is propagated with the ConnectorProfile.

## 5.4.2.3.8 notify\_disconnect

#### Description

This operation notifies a **PortService** of a disconnection between its corresponding port and the other ports. The disconnected connector is identified by the given ID, which was given when the connection was established.

#### **Semantics**

This port invokes the **notify\_disconnect** operation of the next **PortService** in the sequence of the **ConnectorProfile** that was stored when the connection was established. As ports are disconnected, **PortService** references are removed from the **ConnectorProfile**. In this way, the notification of disconnection is propagated by the **notify disconnect** operation.

#### 5.4.2.4 ExecutionContextService

#### Description

An ExecutionContextService exposes an ExecutionContext as an SDO service such that the context may be controlled remotely.

#### **Semantics**

Depending on the implementation, this interface may itself be an execution context (that is, it may be passed to the operations of ComponentAction) or it may represent a remote execution context that is not of type ExecutionContextService.

#### Interface Realizations

• ExecutionContextOperations

#### Generalizations

• SDOService [SDO]

## **Operations**

ExecutionContextService			
no attributes			
operations			
get_profile		ExecutionContextProfile	

## 5.4.2.4.1 get\_profile

#### Description

This operation provides a profile "descriptor" for the execution context.

## 5.4.2.5 FsmService

#### Description

The FsmService interface defines operations necessary for Stimulus Response Processing as an SDO service.

## Generalizations

• SDOService [SDO]

## **Operations**

FsmService				
no attributes				
operations				
get_fsm_profile		FsmProfile		
set_fsm_profile		ReturnCode_t		
	fsm_profile	FsmProfile		

## 5.4.2.5.1 get\_fsm\_profile

## Description

Get the current state of the FSM.

#### **Semantics**

Modifications to the object returned by this operation will not be reflected in the FSM until and unless set\_fsm\_profile is called.

## 5.4.2.5.2 set\_fsm\_profile

## **Description**

This operation will be used to modify the behavior of an FSM as described in Stimulus Response Processing.

# 6 Platform Specific Models

In order to maximize interoperability, this document describes three PSMs that should be considered normative in sections 6.3, 6.4, and 6.5. They correspond to the PSM conformance points outlined in Chapter 2. All PSMs draw on a common set of IDL definitions, which is presented first.

## 6.1 UML-to-IDL Transformation

The PSMs below require IDL definitions for the interfaces, data types, and other model elements from the PIM. They also require IDL representations of the model elements from [UML] on which the PIM depends: Component, Port, Connector, etc. Representing all of the UML in IDL is beyond the scope of this specification. This specification takes a more parsimonious approach.

- IDL definitions for the elements from this specification are provided explicitly in Section 6.2.
- Mapping rules from a subset of UML to IDL are provided in this section. Only those parts of UML that are necessary
  to describe the PIM or critical for the definition of conforming components are described here. Mappings of all other
  UML constructs are implementation-defined.

# 6.1.1 Basic Types and Literals

Primitive types (see Section 5.2.3) shall map to the corresponding IDL primitive types. Specifically:

Character: charDouble: double

Float: floatLong: long

LongDouble: long doubleLongLong: long long

Octet: octetShort: short

• UnsignedLong: unsigned long

• UnsignedLongLong: unsigned long long

• UnsignedShort: unsigned short

WideCharacter: wcharWideString: wstring

The standard UML String and Boolean types are also used by this specification.

Boolean: boolString: string

The literal specifications defined in the PIM (see Section 5.2.4)—as well as those referenced from [UML]—shall be represented as IDL literal values.

## 6.1.2 Classes and Interfaces

UML classes and interfaces shall be represented as IDL interfaces of the same name.

- Each operation or attribute on the UML classifier shall be represented by a corresponding operation or attribute in IDL.
- The general classifiers of UML classes and interfaces shall be represented as inheritance between the corresponding IDL interfaces.

## 6.1.3 Components

RT components shall be represented as IDL components of the same name. As required by the PIM, this component must support the **LightweightRTObject** interface or some subtype thereof, such as **RTObject**.

## 6.1.3.1 Component Inheritance

For each of the UML component's general classifiers that can be represented as a UML interface according to the mapping rules here, the corresponding IDL component shall support the corresponding IDL interface.

If the RTC has a general classifier that is another RTC, the IDL component corresponding to the former shall inherit from that corresponding to the latter.

## 6.1.3.2 Component Ports

Ports providing a single interface shall be represented as facets. Ports requiring a single interface shall be represented as receptacles. Receptacles shall be multiplex unless otherwise constrained by the model. These facets and receptacles shall have the same names as the UML ports.

Ports that expose multiple interfaces, possibly of different polarities, cannot currently be represented in IDL. Therefore, such ports shall be divided into separate facets and receptacles, one per exposed interface, using the following name transformation:

- The name of the port,
- · An underscore,
- The name of the interface's association end.

For example:

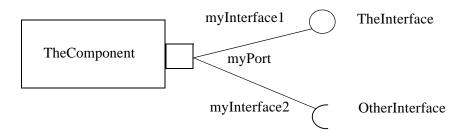


Figure 6.1 - Example of port name mapping

The component in Figure 6.1 would be represented in IDL as follows:

// interfaces TheInterface and OtherInterface previously defined

```
component TheComponent {
    provides TheInterface myPort_myInterface1;
    uses multiple OtherInterface myPort_myInterface2;
};
```

## 6.1.3.3 Composite Components

Implementations may choose to represent composite components—those whose internal parts are interconnected instances of other components—using declarative descriptions of component assemblies (as in [DC]), generated code, or some other means. Deployment and configuration of RTCs is beyond the scope of this specification.

## 6.1.4 Enumerations

Enumerations in the PIM shall be represented as IDL enumerations of the same name. Each attribute shall correspond to a constant within that enumeration.

# 6.1.5 Packages

A UML package shall be represented by an IDL module of the same name.

## 6.2 IDL Definitions

This section defines the representations of the PIM elements in IDL. These definitions have been created using the rules from Section 6.1 as extended below. Several minor changes have been made from the PIM in order to make the appearance of the API more natural in IDL (and in programming language APIs generated from it).

The IDL definitions themselves may be found in an annex to this document as well as in the accompanying file RTC.idl (see <a href="http://www.omg.org/spec/RTC/20070801">http://www.omg.org/spec/RTC/20070801</a>). Only the latter is normative.

## 6.2.1 Classes and Interfaces

Because UML classes and interfaces are both represented by IDL interfaces (see Section 6.1.2), the types **ExecutionContext** and **ExecutionContextOperations** are essentially identical. Therefore, in the PSMs, they have been collapsed to a single IDL interface "**ExecutionContext**."

## 6.2.2 Stereotypes

Stereotypes have been represented as IDL interfaces intended to be supported by components having those stereotypes. This convention provides a convenient way to add the necessary operations to supporting components. Two additional rules apply:

- Stereotypes that constrain their extended types to realize a given interface have been unified with that interface.
- IDL interfaces representing UML stereotypes are named with an initial capital letter to reflect common naming conventions for IDL interfaces.

#### Specifically:

- lightweightRTComponent has been unified with the LightweightRTObject interface.
- dataFlowComponent is represented by the DataFlowComponent interface.
- fsm is represented by the Fsm interface.
- fsmParticipant is represented by the FsmParticipant interface.
- multiModeComponent has been unified with the MultiModeComponentObject interface.
- rtComponent has been unified with the RTObject interface.

## 6.2.3 Return Codes

The PIM allows ReturnCode\_t to be modeled in a PSM as either an actual return value or as an exception. The IDL definition chooses a return value as the lowest common denominator mapping across programming languages, and therefore retains the enumeration representation from Section 6.1.4. However, in an annex and RTC.idl, since a word confliction is expected in mapping to other languages, OK and ERROR are replaced to RTC\_OK and RTC\_ERROR. This convention should not be construed to overrule the more general allowance in the PIM.

# 6.2.4 Packages

In the PIM, the RTC package has sub-packages. These sub-packages have been flattened in the PSMs in order to simplify the API in those language bindings (such as C) that reflect IDL modules in contained identifier names.

## 6.3 Local PSM

RTCs implemented using this PSM communicate within a single local application without intervening network communication; RTCs, assembly connectors, and delegation connectors are implemented as local objects. Specifically, the presence of a CORBA ORB shall not be required.

This communication model is especially important for RTCs with very strict timeliness, reliability, and determinism requirements. It is also applicable to resource-constrained environments. (The default communication styles in Simulink from The MathWorks, LabView from National Instruments, and Constellation from RTI are analogous to this PSM.)

## 6.3.1 IDL Transformation Rules

The programming interfaces for this PSM are implied by the IDL and IDL mapping rules from Section 6.1 and Section 6.2 as extended below. One additional exception applies: because this PSM is not CORBA-based, any CORBA-specific base types supplied by the IDL-to-programming language mapping shall be omitted from any generated code.

This PSM relies on IDL type definitions to provide a programming language-agnostic programming interface to an RTC-based middleware. Therefore, it is only well-defined in the context of a particular IDL-to-programming language mapping. To derive the RTC API in the target programming language, that mapping shall be applied to the input IDL with the additional restrictions noted in this section. These restrictions stem from the need to eliminate usages of CORBA-specific types that may be incidentally implied by a particular IDL-to-programming language mapping.

The restrictions in this section currently apply to the C++ language only, and specifically to the OMG-standard IDL-to-C++ mapping, although future revisions may introduce explicit support for additional programming language mappings. Until that time, conforming implementations of this PSM with other programming language mappings will not be possible.

#### 6.3.1.1 Mapping for Interfaces and Reference Types

#### 6.3.1.1.1 Implicit Local Interfaces

Since object distribution is beyond the scope of this PSM, all IDL interfaces shall be treated as local interfaces with respect to code generation. That is, any code generated shall be as if the interface declaration was prefixed with the IDL local keyword.

## 6.3.1.1.2 CORBA-Specific Base Interfaces

The class generated for an IDL interface shall not extend any CORBA-specific base class such as CORBA::Object or CORBA::LocalObject. The programming interfaces provided by those base classes are therefore non-normative with respect to this PSM.

- Nil values shall be represented by the value **NULL**, eliminating the need for **is nil** and similar operations.
- Object references are represented as simple pointers (see Section 6.3.1.1.4), and can therefore be compared directly, eliminating the need for CORBA::Object::is equivalent.
- Runtime type identification is beyond the scope of this PSM. Implementations that require it must either rely on the support built into the target programming language or provide an alternative implementation-specific mechanism.

#### 6.3.1.1.3 Abstract Interfaces

This specification does not use any abstract interfaces, nor does any specification on which it depends. Furthermore, abstract interfaces do not result from any of the UML-to-IDL transformation rules in this specification. Therefore, the IDL mapping of abstract interfaces is outside the scope of this PSM.

#### 6.3.1.1.4 Pointer Types

Support for the "smart pointer" \_ptr, \_var, and \_out types shall not be required. Instead, object references shall be represented with native pointers. Memory management is the responsibility of the party that allocates it: either the middleware or the component. This convention applies not only to interface types, but also to strings, wide strings, and other non-basic types.

Because pointers are used directly, the \_duplicate operation on interface types is not needed and is therefore not normative with respect to this PSM.

## 6.3.1.2 Mapping for Basic Data Types

IDL basic data types shall map to type definitions in the RTC namespace having the same names as the corresponding PIM types:

• short: RTC::Short

• long: RTC::Long

long long:
 RTC::LongLong

• unsigned short: RTC::UnsignedShort

unsigned long: RTC::UnsignedLong

unsigned long long: RTC::UnsignedLongLong

float: RTC::Float

• double: RTC::Double

long double: RTC::LongDouble

char: RTC::Char

• wchar: RTC::WideChar

• boolean: RTC::Boolean

• octet: RTC::Octet

Programmers concerned with portability should use the **RTC** types. However, some may feel that using these types with the **RTC** qualification impairs readability. As in the OMG-standard IDL-to-C++ mapping, on platforms where the C++ data type is guaranteed to be identical to the OMG IDL data type, a compliant implementation may generate the native C++ type.

## 6.3.1.3 Mapping for String Types

IDL strings map to **char\*** as in the OMG-standard IDL-to-C++ mapping. For dynamic allocation of strings, compliant implementations must use the following functions from the RTC namespace:

```
// C++
namespace RTC {
    char* string_alloc(RTC::UnsignedLong len);
    char* string dup(const char*);
```

```
void string_free(char*);
}
```

The behavior of these functions shall be the same as the equivalent functions defined in the OMG-standard IDL-to-C++ mapping specification.

## 6.3.1.4 Mapping for Wide String Types

IDL wide strings, whether bounded or unbounded, map to RTC::WideChar\*. For dynamic allocation of wide strings, compliant implementations must use the following functions from the RTC namespace:

```
// C++
namespace RTC {
    RTC::WideChar* wstring_alloc(RTC::UnsignedLong len);
    RTC::WideChar* wstring_dup(const RTC::WideChar*);
    void wstring_free(RTC::WideChar*);
}
```

The behavior of these functions shall be the same as the equivalent functions defined in the OMG-standard IDL-to-C++ mapping specification.

## 6.3.1.5 Mapping for Fixed Types

A mapping for the IDL fixed point type is outside of the scope of this PSM.

## 6.3.1.6 Mapping for Any Type

The IDL any type shall be represented in generated code as an instance of the type RTC::Any. The programming interfaces and semantics are as described in the OMG-standard IDL-to-C++ mapping; only the namespace is changed.

As described elsewhere, the types CORBA::AbstractBase, CORBA::ValueBase, CORBA::TypeCode, and CORBA::Fixed are non-normative with respect to this PSM. Therefore, the Any member structures and operations that refer to these types are also non-normative.

## 6.3.1.7 Mapping for Valuetypes

IDL **valuetype**s are used neither by this specification nor by the SDO specification. Furthermore, they do not result from any of the UML-to-IDL mappings defined for this specification. Therefore, the IDL-to-programming language mappings for **valuetype**s are beyond the scope of this PSM.

#### 6.3.1.8 Exceptions

Operations defined in this specification report errors by means of ReturnCode\_t objects, not IDL exception objects. They shall therefore refrain from throwing system exceptions and shall instead fail with the appropriate ReturnCode\_t result, such as BAD PARAMETER, OUT OF RESOURCES, or simply ERROR.

However, the SDO specification does rely on exceptions for error reporting. Additionally, this PSM specifies that certain programming interfaces defined by the OMG-standard IDL-to-C++ mapping (identified elsewhere in this PSM) shall be provided in the RTC namespace. Unless otherwise noted, these APIs do not use ReturnCode\_t to report errors. Therefore, the types Exception, SystemException, CompletionStatus, and UserException shall be

defined in the RTC namespace. These shall take the place of their CORBA namespace equivalents in all programming interfaces, including as base classes for user-defined and standard exception types. They shall function identically to their corresponding types from the CORBA namespace. Standard subclasses of **SystemException** shall also exist in the RTC namespace rather than in the **CORBA** namespace.

The OMG-standard IDL-to-C++ mapping allows for C++ implementations that do not support exception handling. This allowance is made through the provision of an **Environment** pseudo-interface. Implementations of this PSM that do not support exceptions shall provide a class RTC::Exception that is functionally equivalent to CORBA::Environment.

## 6.3.1.9 Mapping for Components

The mapping for IDL component-related concepts relies on the "equivalent interfaces" outlined in the CCM specification. These IDL-to-IDL transformations remain normative.

The class generated for an IDL component shall not extend any CORBA- or CCM-specific base class such as Components::CCMObject. The programming interfaces provided by that base class, and its base classes, are therefore non-normative with respect to this PSM.

#### 6.3.1.9.1 Mapping for Receptacles

## **Simplex Receptacles**

A uses declaration of the following form:

```
uses <interface type> <receptacle name>;
```

results in the following equivalent operations defined in the component interface.

```
ReturnCode_t connect_<receptacle_name>(
    in <interface_type> conxn);
<interface_type> disconnect_<receptacle_name>();
<interface_type> get_connection_<receptacle_name>();
```

These definitions match those defined by the Lightweight CCM specification except with respect to error reporting. They report error conditions using the **ReturnCode\_t** mechanism defined in this specification rather than with explicit exceptions. The disconnect operation shall indicate failure with a nil return result.

#### **Multiplex Receptacles**

A uses declaration of the following form:

```
uses multiple <interface_type> <receptacle_name>;
```

results in the following equivalent operations defined in the component interface.

```
struct <receptacle_name>Connection {
  <interface_type> objref;
  RTC::Cookie ck;
};
```

sequence<<receptacle\_name>Connection> <receptacle\_name>Connections;

```
RTC::Cookie connect_<receptacle_name>(
    in <interface_type> connection);
<interface_type> disconnect_<receptacle_name>(
    in RTC::Cookie ck);
<receptacle_name>Connections get_connections_<receptacle_name>();
```

The above declarations are similar to those defined by the Lightweight CCM specification with two exceptions: error handling and the replacement of Components::Cookie with RTC::Cookie. The latter shall be identical to the former apart from the change in namespace.

The disconnect operation shall report errors by means of a nil return result.

## 6.3.2 Interfaces

Existing modeling and development tools used by robotics developers, such as Simulink from The MathWorks and LabView from National Instruments, often rely on data-only connections between program modules. The connection types in such cases are not interfaces in the IDL sense; they may be primitives, arrays, or simple structures. This approach has several benefits from an implementation standpoint:

- Robotics developers, who are often not software engineers, think in data-centric terms, so this ability is familiar and comfortable for them.
- Using unencapsulated data decreases the number of function calls necessary along the critical path of a high-rate realtime application.
- Legacy robotics code is frequently not interface-centric, and integrating with it is simpler if legacy types can be used directly.

For these reasons, UML interfaces may be represented by any data type that can be expressed in the target language as long as it fulfills the contract described by the interface.

#### 6.3.3 Connectors

Connectors shall be represented as instances of the type of the connected ports. This type may contain data, methods, or both. If the type of the connector contains data, implementations must ensure that all ports connected to it observe the same values at all times.

If the type of the connector contains methods, the connector implementation is responsible for performing any multiplexing or de-multiplexing of method calls that may be necessary when a connector is not one-to-one. When a single caller (i.e., an operation defined by a required interface) is connected to multiple callees (i.e., the same operation within multiple instances of the matching provided interface), the implementation must invoke all callee methods. However, the handling of their return results or exceptions is implementation-defined.

#### 6.3.4 SDO

The Introspection package (Section 5.4) of the PIM relies on the Super Distributed Objects specification [SDO]. Implementations that support that package shall use the IDL provided with that specification's CORBA PSM as transformed according to this PSM.

# 6.4 Lightweight CCM PSM

In this PSM, RTCs are mapped to CCM components supporting the relevant IDL interfaces described in Section 6.1 and Section 6.2.

## 6.4.1 Connectors

CCM (including Lightweight CCM) does not currently support explicit connector objects. Therefore, connectors are implicitly defined by the connections between components.

#### 6.4.2 SDO

The Introspection package (Section 5.4) of the PIM relies on the Super Distributed Objects specification [SDO]. Implementations that support that package shall use that specification's CORBA PSM.

## 6.4.3 Behavior

This specification's PIM defines behavioral constraints that conforming implementations must respect. (For example, callbacks of data flow participant components must be invoked in a well-defined order, as described in Section 5.3.1). A middleware that hosts RT components and is implemented using CCM may choose how to satisfy those constraints. However, some non-normative suggestions are provided here.

- The execution semantics in this specification can be expressed in part using the Execution Models defined by the Streams for CCM Specification ([CCM STREAM], section 10.2).
- The implementation of a distributed execution context (in which RTCs on different nodes participate in the same context) may benefit from the *distributable thread* feature of the Real-Time CORBA specification [RT CORBA].

## 6.5 CORBA PSM

In this PSM, RTCs are mapped to CORBA interfaces extending the relevant IDL interfaces described in Section 6.1 and Section 6.2.

# 6.5.1 Mapping for Components

The mapping for IDL component-related concepts relies on the "equivalent interfaces" outlined in [CCM]. These IDL-to-IDL transformations remain normative.

The class generated for an IDL component shall not extend the CCM-specific base class Components::CCMObject. Instead, the class shall extend CORBA::Object as is typical of CORBA interfaces.

#### 6.5.1.1 Mapping for Receptacles

#### 6.5.1.1.1 Simplex Receptacles

A uses declaration of the following form:

uses <interface\_type> <receptacle\_name>;

results in the following equivalent operations defined in the component interface.

```
ReturnCode_t connect_<receptacle_name>(
    in <interface_type> conxn);
<interface_type> disconnect_<receptacle_name>();
<interface_type> get_connection_<receptacle_name>();
```

These definitions match those defined by the Lightweight CCM specification except with respect to error reporting. They report error conditions using the ReturnCode\_t mechanism defined in this specification rather than with explicit exceptions. The disconnect operation shall indicate failure with a nil return result.

## 6.5.1.1.2 Multiplex Receptacles

A uses declaration of the following form:

```
uses multiple <interface_type> <receptacle_name>;
```

results in the following equivalent operations defined in the component interface.

```
struct <receptacle_name>Connection {
  <interface_type> objref;
  RTC::Cookie ck;
};
```

sequence<<receptacle\_name>Connection> <receptacle\_name>Connections;

```
RTC::Cookie connect_<receptacle_name>(
    in <interface_type> connection);
<interface_type> disconnect_<receptacle_name>(
    in RTC::Cookie ck);
<receptacle_name>Connections get_connections_<receptacle_name>();
```

The above declarations are similar to those defined by the Lightweight CCM specification with two exceptions: error handling and the replacement of Components::Cookie with RTC::Cookie. The latter shall be identical to the former apart from the change in namespace.

The disconnect operation shall report errors by means of a nil return result.

# 6.5.2 Mapping for Connectors

CORBA does not currently support explicit connector objects. Therefore, connectors are implicitly defined by the references between objects.

#### 6.5.3 SDO

The Introspection package (Section 5.4) of the PIM relies on the Super Distributed Objects specification [SDO]. Implementations that support that package shall use that specification's CORBA PSM.

## 6.5.4 Behavior

This specification's PIM defines behavioral constraints that conforming implementations must respect. (For example, callbacks of data flow participant components must be invoked in a well-defined order, as described in Section 5.3.1). A middleware that hosts RT components and is implemented using CORBA may choose how to satisfy those constraints. However, some non-normative suggestions are provided here.

• The implementation of a distributed execution context (in which RTCs on different nodes participate in the same context) may benefit from the distributable thread feature of the Real-Time CORBA specification [RT CORBA].

# Annex A RTC IDL

## (non-normative)

The normative IDL used by the PSMs in this specification is contained in the file RTC.idl that accompanies this document. A non-normative copy of that file's contents is provided here for convenience.

```
// RTC.idl
```

```
#include "SDOPackage.idl"
#pragma prefix "omg.org"
#define EXECUTION HANDLE TYPE NATIVE long
module RTC {
   typedef EXECUTION_HANDLE_TYPE_NATIVEExecutionContextHandle_t;
                                            Uniqueldentifier;
   typedef SDOPackage::UniqueIdentifier
   typedef SDOPackage::NVList
                                            NVList;
   enum ReturnCode t {
      RTC OK,
      RTC ERROR,
      BAD PARAMETER,
      UNSUPPORTED,
      OUT OF RESOURCES,
      PRECONDITION NOT MET
   };
   enum LifeCycleState {
      CREATED STATE,
      INACTIVE STATE,
      ACTIVE_STATE,
      ERROR STATE
   };
   interface ExecutionContext;
   typedef sequence<ExecutionContext> ExecutionContextList;
   interface ComponentAction {
      ReturnCode t on initialize();
      ReturnCode t on finalize();
      ReturnCode t on startup(
          in ExecutionContextHandle_t exec_handle);
      ReturnCode t on shutdown(
          in ExecutionContextHandle_t exec_handle);
      ReturnCode t on activated(
          in ExecutionContextHandle t exec handle);
```

ReturnCode t on deactivated(

```
in ExecutionContextHandle t exec handle);
   ReturnCode t on aborting(
       in ExecutionContextHandle t exec handle);
   ReturnCode t on error(in ExecutionContextHandle t exec handle);
   ReturnCode_t on_reset(in ExecutionContextHandle_t exec_handle);
};
interface LightweightRTObject : ComponentAction {
   ReturnCode t initialize();
   ReturnCode t finalize();
   boolean is_alive(in ExecutionContext exec_context);
   ReturnCode t reset();
   ReturnCode t exit();
   ExecutionContextHandle t attach context(
       in ExecutionContext exec context);
   ReturnCode_t detach_context(
       in ExecutionContextHandle t exec handle);
   ExecutionContext get context(
       in ExecutionContextHandle t exec handle):
   ExecutionContextList get_owned_contexts();
   ExecutionContextList get_participating_contexts();
};
enum ExecutionKind {
   PERIODIC,
   EVENT DRIVEN,
   OTHER
};
interface ExecutionContext {
   boolean is_running();
   ReturnCode t start();
   ReturnCode_t stop();
   double get rate();
   ReturnCode t set rate(in double rate);
   ReturnCode_t add_component(
       in LightweightRTObject comp);
   ReturnCode_t remove_component(
       in LightweightRTObject comp);
   ReturnCode t activate component(
       in LightweightRTObject comp);
   ReturnCode t deactivate component(
       in LightweightRTObject comp);
   ReturnCode_t reset_component(
       in LightweightRTObject comp);
   LifeCycleState get component state(
       in LightweightRTObject comp);
   ExecutionKind get_kind();
};
interface DataFlowComponentAction {
```

```
ReturnCode t on execute(
       in ExecutionContextHandle t exec handle);
   ReturnCode_t on_state_update(
       in ExecutionContextHandle t exec handle);
   ReturnCode_t on_rate_changed(
       in ExecutionContextHandle t exec handle);
};
interface DataFlowComponent : DataFlowComponentAction {
interface Fsm {
};
interface FsmParticipantAction {
   ReturnCode_t on_action(
       in ExecutionContextHandle t exec handle);
};
interface FsmParticipant : FsmParticipantAction {
interface Mode {
};
interface ModeCapable {
   Mode get_default_mode();
   Mode get_current_mode();
   Mode get current mode in context(
       in ExecutionContext exec context);
   Mode get_pending_mode();
   Mode get_pending_mode_in_context(
       in ExecutionContext exec_context);
   ReturnCode t set mode(
       in Mode new mode,
       in boolean immediate);
};
interface MultiModeComponentAction {
   ReturnCode t on mode changed(
       in ExecutionContextHandle t exec handle);
};
interface MultiModeObject : ModeCapable,
                          MultiModeComponentAction {
};
interface RTObject;
enum PortInterfacePolarity {
   PROVIDED,
```

```
REQUIRED
};
struct PortInterfaceProfile {
   string instance_name;
   string type name;
   PortInterfacePolarity polarity;
};
typedef sequence<PortInterfaceProfile> PortInterfaceProfileList;
interface PortService;
typedef sequence<PortService> PortServiceList;
typedef sequence<RTObject> RTCList;
struct ConnectorProfile {
   string name;
   Uniqueldentifier connector id;
   PortServiceList ports;
   NVList properties;
};
typedef sequence<ConnectorProfile> ConnectorProfileList;
struct PortProfile {
   string name;
   PortInterfaceProfileList interfaces;
   PortService port ref;
   ConnectorProfileList connector profiles;
   RTObject owner;
   NVList properties;
};
typedef sequence<PortProfile> PortProfileList;
struct ExecutionContextProfile {
   ExecutionKind kind;
   double rate;
   RTObject owner;
   RTCList participants;
   NVList properties;
};
typedef sequence<ExecutionContextProfile>
ExecutionContextProfileList;
interface FsmObject {
   ReturnCode_t send_stimulus(
       in string message,
       in ExecutionContextHandle t exec handle);
};
```

```
struct FsmBehaviorProfile {
   FsmParticipantAction action_component;
   Uniqueldentifier id;
};
typedef sequence<FsmBehaviorProfile> FsmBehaviorProfileList;
struct FsmProfile {
   FsmBehaviorProfileList behavior_profiles;
};
interface FsmService : SDOPackage::SDOService {
    FsmProfile get fsm profile();
    ReturnCode_t set_fsm_profile(in FsmProfile fsm_profile);
};
struct ComponentProfile {
   string instance name;
   string type name;
   string description;
   string version;
   string vendor;
   string category;
   PortProfileList port profiles;
   RTObject parent;
   NVList properties;
};
typedef sequence<ComponentProfile> ComponentProfileList;
interface PortService : SDOPackage::SDOService {
   PortProfile get_port_profile();
   ConnectorProfileList get_connector_profiles();
   ConnectorProfile get connector profile(
       in Uniqueldentifier connector id);
   ReturnCode t connect(
       inout ConnectorProfile connector profile);
   ReturnCode_t disconnect(in UniqueIdentifier connector_id);
   ReturnCode t disconnect all();
   ReturnCode t notify connect(
       inout ConnectorProfile connector_profile);
   ReturnCode t notify disconnect(
       in UniqueIdentifier connector_id);
};
interface ExecutionContextService: ExecutionContext,
                                     SDOPackage::SDOService {
   ExecutionContextProfile get_profile();
};
typedef sequence<ExecutionContextService>
```

## **ExecutionContextServiceList**;

```
interface RTObject : LightweightRTObject, SDOPackage::SDO {
     ComponentProfile get_component_profile();
     PortServiceList get_ports();
    };
};
```

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